

the right majors, minors & concentrations
education?

for students' academic and career success
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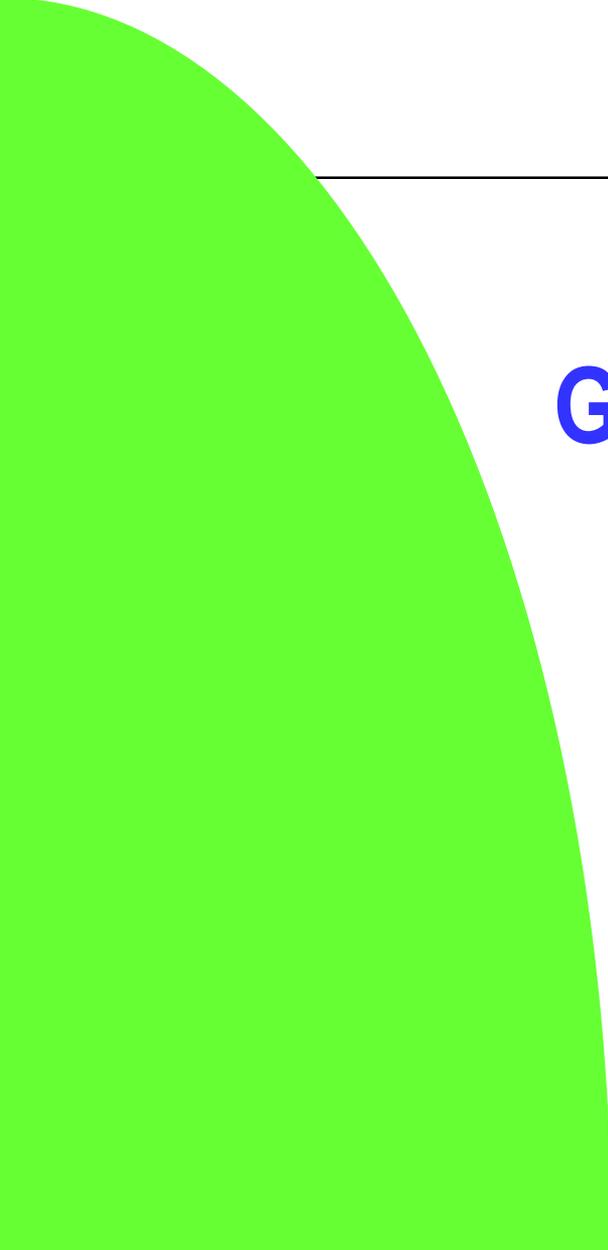
Graphs and Graph Problems

Y. PARK • DEPT. OF IS&IS, BRUNEL UNIVERSITY

1/7

Prof. Young Park





Graphs as Non-linear Data Structures

What Is a Graph?

- A **nonlinear** data structure consisting of a set V of nodes (vertices) and a set E of links between the nodes (edges or arcs).
 - ➡ $G = (V, E)$
 - ➡ $V = \{v_1, v_2, \dots, v_n\}$
 - ➡ $E = \{ (v_i, v_j) \}$
- An edge can connect any pair of vertices.
 - ➡ v is **adjacent** to u when (u, v) in E .
- **Like lists & trees, but more general!**

What Graphs?

- Undirected graphs
- Directed graphs
- Acyclic graphs
- Connected graphs
- Labeled (Weighted) graphs
- Sparse/Dense graphs
- ...

Paths in Graphs

- A **path** in a graph is
 - An alternating sequence of vertices and edges in which all edges and vertices are distinct.
 - A sequence of vertices in which all vertices are distinct.
 - ☞ $P = \langle v_1, v_2, \dots, v_k \rangle$ where $(v_i, v_j) \in E$
- The length of a path is
 - The number of edges on the path.

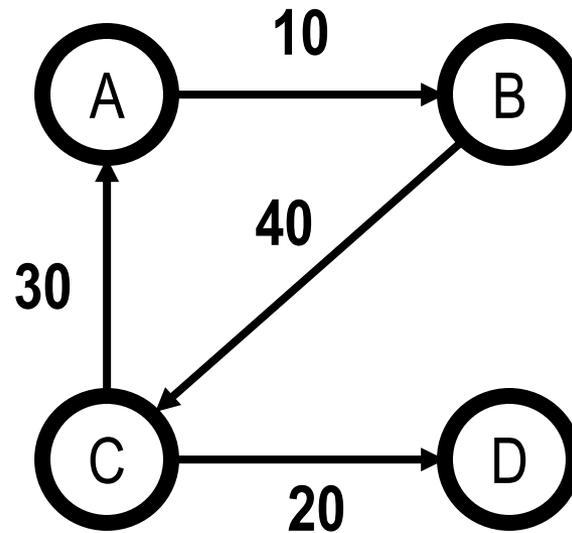
Cycles in Graphs

- A **cycle** in a graph is
 - A path in which the first vertex and the last vertex is the same.

Labeled Graphs

- A graph $G = (V, E)$ is called a **labeled graph**
 - If its edges are assigned data of one kind of another.
 - ☞ labeled edges
 - If its vertices are assigned data of one kind of another.
 - ☞ labeled vertices
 - If its edges and vertices are assigned data of one kind of another.
 - ☞ labeled edges and vertices
 - The assigned some numerical value is called the weight or cost. (called **Weighted Graphs**)

Example: Edge-Weighted Graphs



Sparse Graphs vs Dense Graphs

- A graph $G = (V, E)$ is said to be **sparse**
 - If it has relatively few edges.
 - $|E| = O(|V|)$ edges
- A graph $G = (V, E)$ is said to be **dense**
 - If it has many edges.
 - $|E| = O(|V|^2)$ edges

Why Graphs in Problem Solving?

- Many real-world problems can be reduced to problems on graphs.
 - Often a problem can be represented as a graph and the solution to the problem is obtained by solving a problem on the corresponding graph!
 - Many problems can be solved by asking an appropriate question about paths in a graph!

Graphs vs Other Data Structures

- A graph is a **generalized** data structure of other data structures
 - Lists
 - Trees
 - Heaps
 - ...

► QUIZ?

- Graphs vs. Trees?
- Dense graphs vs. Sparse graphs?



Representation of Graphs

How to Represent Graphs?

- An array-based representation
 - Use an **adjacency matrix**
- A pointer-based representation
 - Use an **adjacency list**

1. An Adjacency Matrix-based Representation

- Use a two-dimensional array - matrix
- For an **undirected** graph $G = (V, E)$
 - $A[i][j] =$
 - 1 if (v_i, v_j) in E
 - 0 otherwise
- For an edge-weighted **undirected** graph $G = (V, E)$
 - $A[i][j] =$
 - weight (numeric label) if (v_i, v_j) in E
 - ∞ otherwise
- $A[i][j] = A[j][i]$

An Adjacency Matrix-based Representation

- For a **directed** graph $G = (V, E)$

→ $A[i][j] =$

→ 1 if (v_i, v_j) in E

→ 0 otherwise

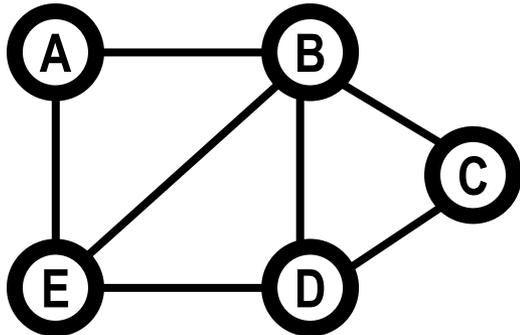
- For an edge-weighted **directed** graph $G = (V, E)$

→ $A[i][j] =$

→ weight (numeric label) if (v_i, v_j) in E

→ ∞ otherwise

Example: An Adjacency Matrix-based Representation



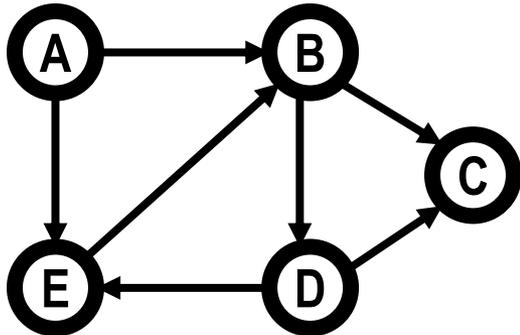
V

0	A
1	B
2	C
3	D
4	E
·	·
·	·
·	

E

	0	1	2	3	4	...
0	0	1	0	0	1	
1	1	0	1	1	1	
2	0	1	0	1	0	
3	0	1	1	0	1	
4	1	1	0	1	0	
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Example: An Adjacency Matrix-based Representation



V

0	A
1	B
2	C
3	D
4	E
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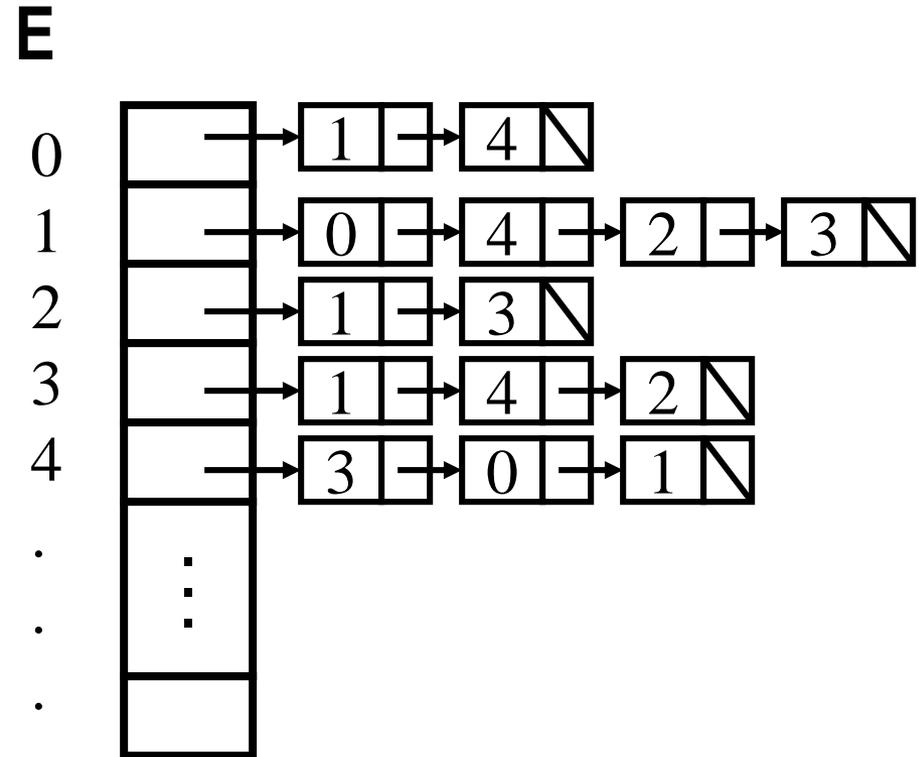
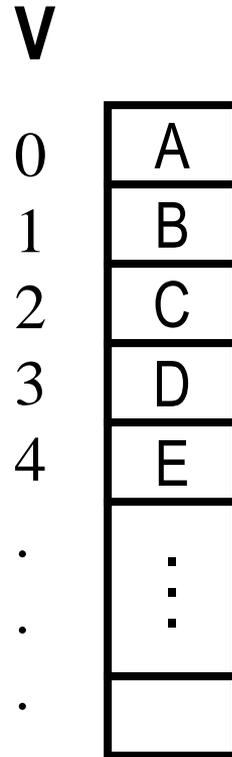
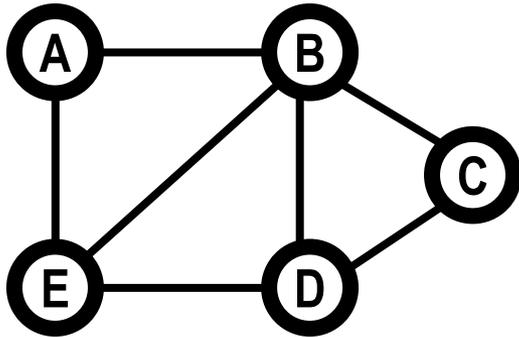
E

	0	1	2	3	4	...
0	0	1	0	0	1	
1	0	0	1	1	0	
2	0	0	0	0	0	
3	0	0	1	0	1	
4	0	1	0	0	0	
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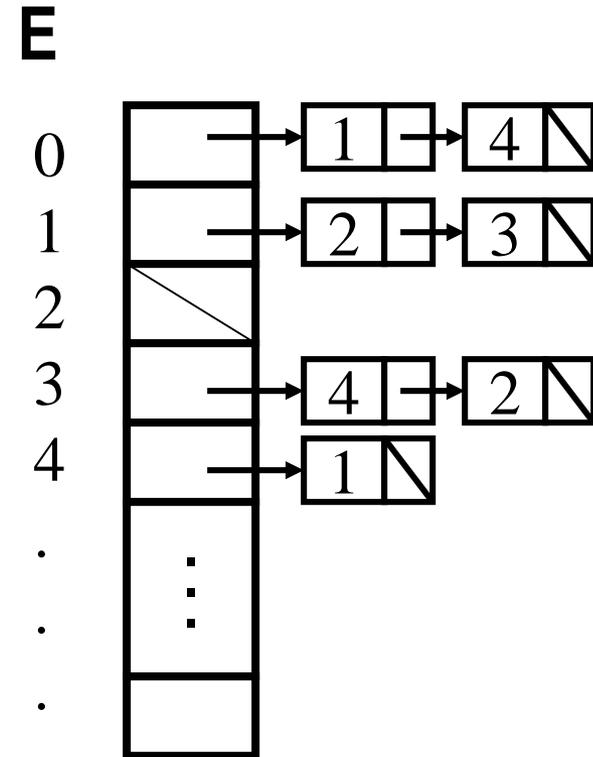
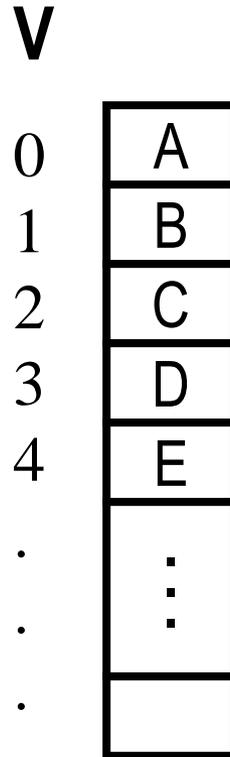
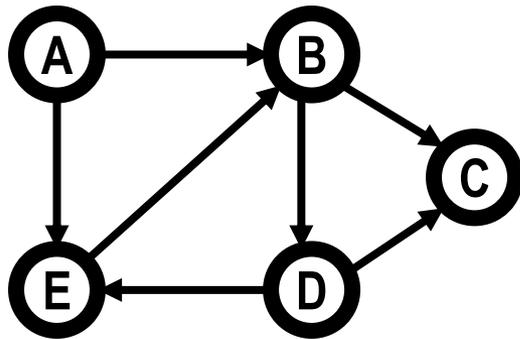
2. An Adjacency (linked) List-based Representation

- Use a linked-list.
- For an **undirected or directed** graph $G = (V, E)$
 - Use $|V|$ linked lists - one for each vertex.
 - The linked list for vertex v_i in V contains the set of vertices adjacent to v_i .
- For an edge-weighted **undirected or directed** graph $G = (V, E)$
 - Use $|V|$ linked lists - one for each vertex.
 - The linked list for vertex v_i in V contains the set of vertices adjacent to v_i and weights (numeric labels).

Example: An Adjacency (linked) List-based Representation



Example: An Adjacency (linked) List-based Representation



Comparison of Graph Representations

- Space comparison
- Time comparison

1. Space Comparison of Graph Representations

- Adjacent matrix
 - $O(|V|^2)$ space
- Adjacent list
 - $O(|V| + |E|)$ space

2. Time Comparison of Graph Representations

- Operation: **Adding or removing an edge.**
 - Adjacent matrix
 - 👉 $O(1)$
 - 👉 $O(1)$
 - Adjacent list
 - 👉 $O(1)$
 - 👉 $O(|E|)$

Time Comparison of Graph Representations

- Operation: **Determine whether there is an edge from vertex u to another vertex v .**
 - Adjacent matrix
 - ☞ $O(1)$ time
 - Adjacent list
 - ☞ $O(|E|)$ time

Time Comparison of Graph Representations

- Operation: **Find all edges.**

- Adjacent matrix

- ☞ $O(|V|^2)$ time

- Adjacent list

- ☞ $O(|V|+|E|)$ time

Time Comparison of Graph Representations

- Operation: **Find all vertices v that are adjacent to a given vertex u s.t. (u,v) .**
 - Adjacent matrix
 - ☞ $O(|V|)$
 - Adjacent list
 - ☞ $O(|E|)$

Time Comparison of Graph Representations

- Operation: **Find all vertices u that a given vertex v is adjacent to (u,v) .**
 - Adjacent matrix
 - ☞ $O(|V|)$ time
 - Adjacent list
 - ☞ $O(|V|+|E|)$ time

► QUIZ?

- Adjacent matrix vs. Adjacent list?



Operations on Graphs

Graphs Traversals

- Often it is useful to visit the vertices of a graph in some specific order.
- Two common ways:
 - **Depth-first search/traversal**
 - **Breadth-first search/traversal**

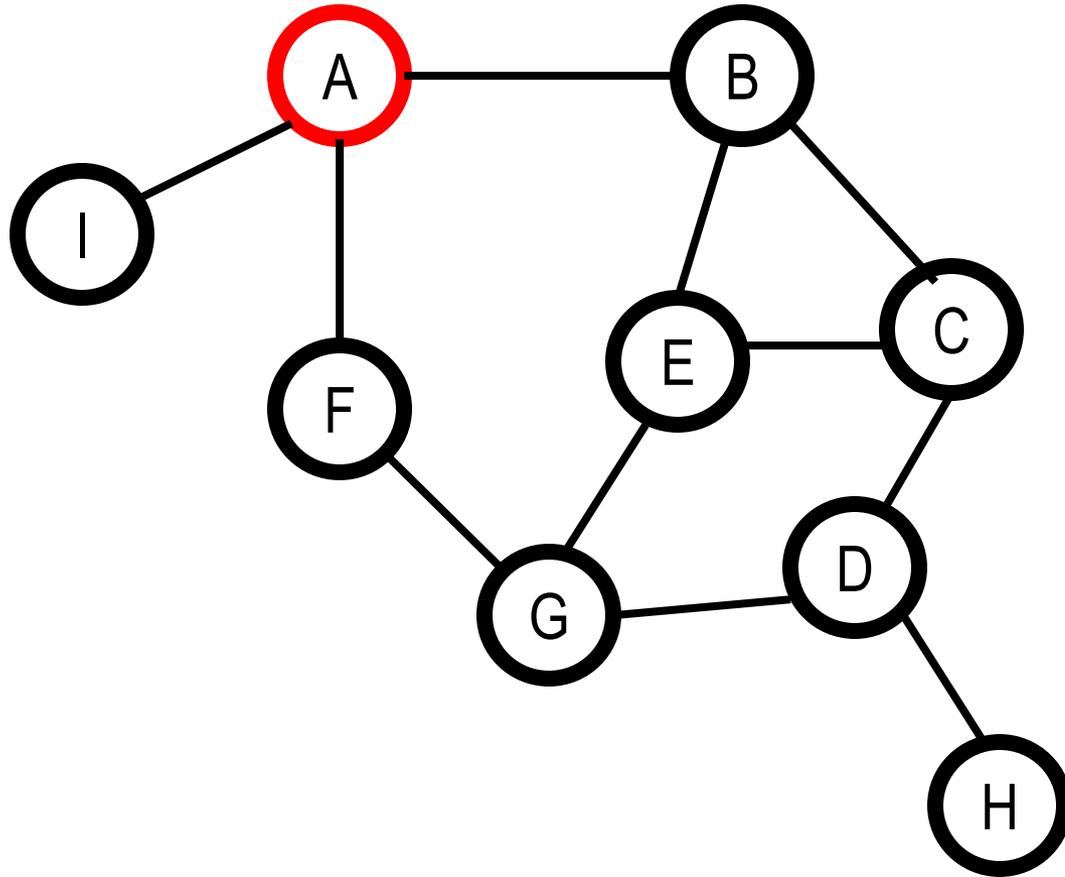
Tree Traversal Vs Graph Traversal?

- Tree
 - Acyclic
- Graphs
 - Could have cycles!
 - Unlike tree traversal, need to specifically guard against repeating a path from a cycle.
 - Mark each vertex as “visited” when we encounter it and not consider visited vertices more than once.

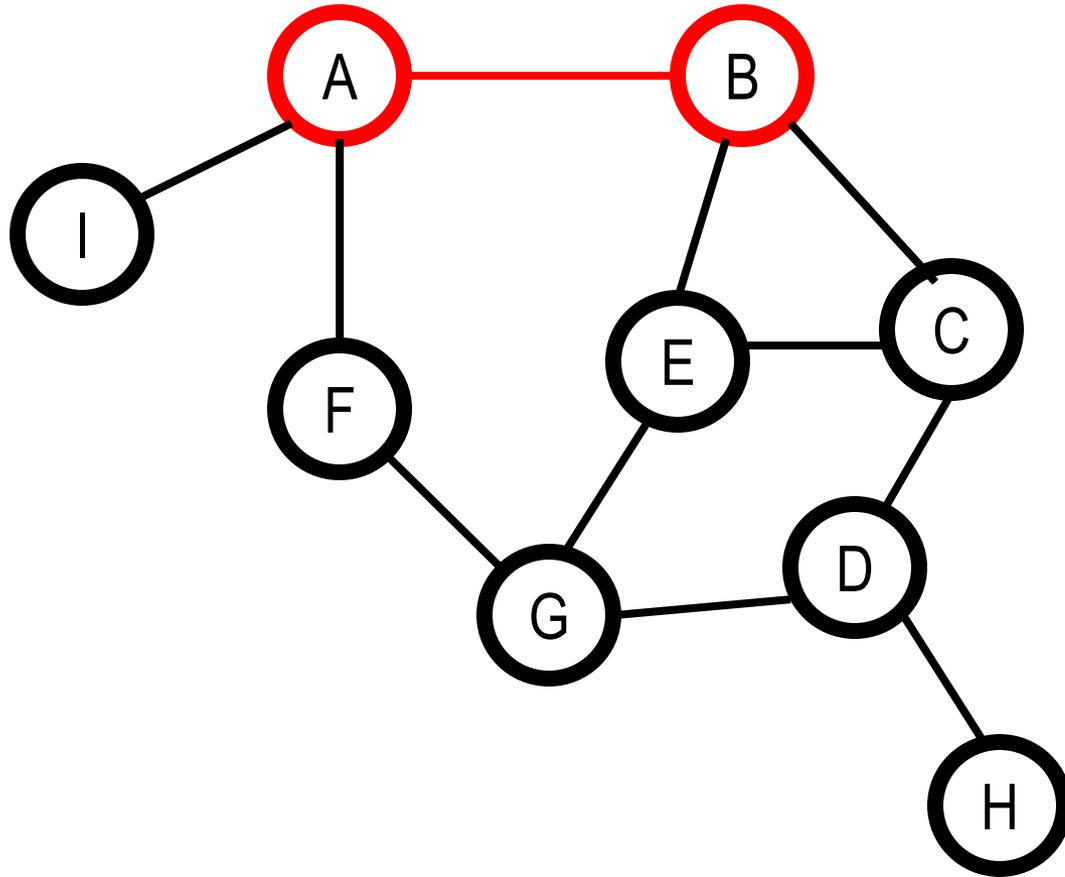
Depth-First Search (DFS) - Recursive

- DFS(u):
 - Mark u as visited;
 - for each **unvisited** vertex v **adjacent** to u (i.e., edge (u,v)) do
 - ☞ DFS(v)

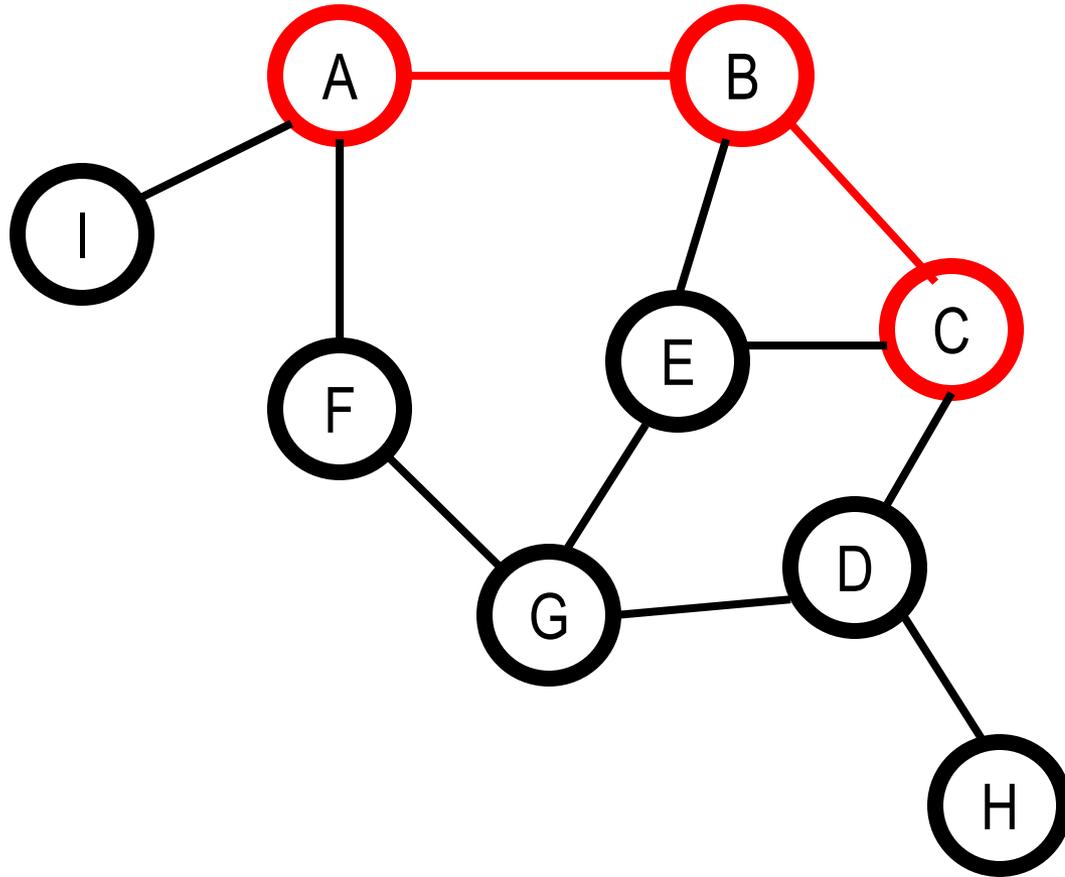
Example: DFS(A) - Recursive - of a connected, undirected graph



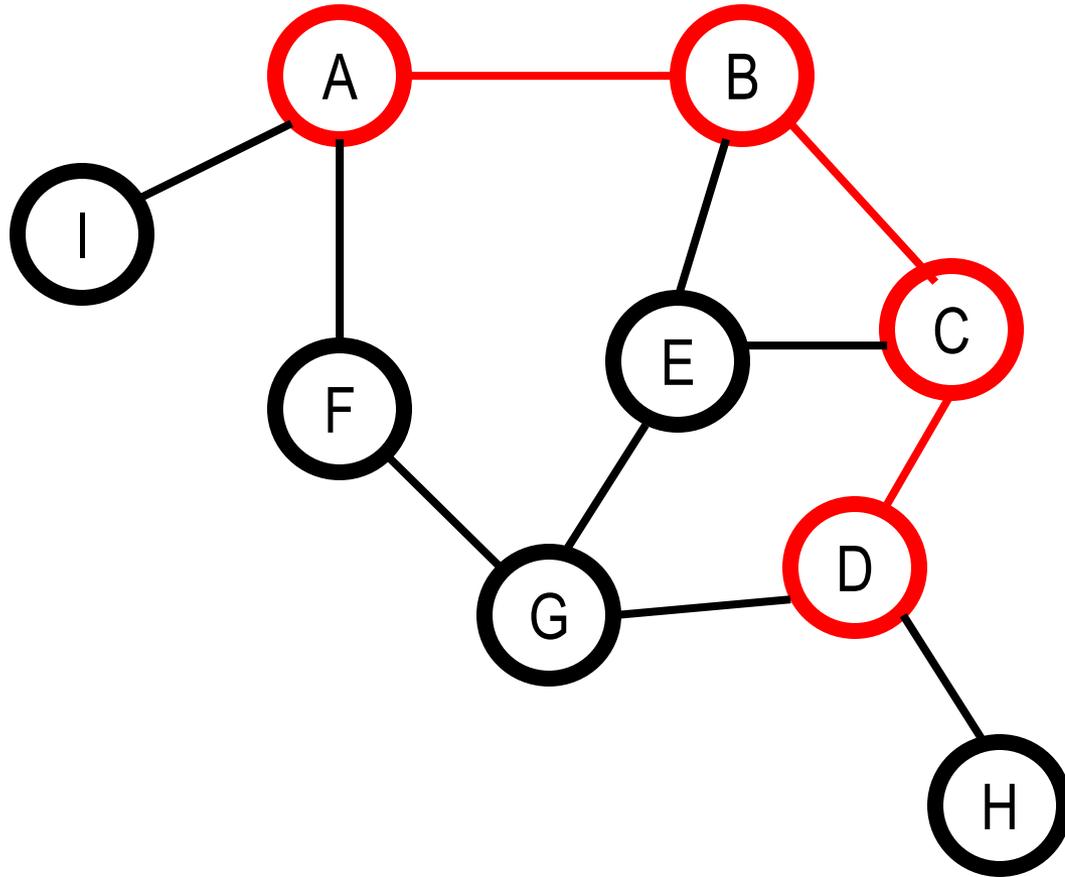
Example: DFS(A)



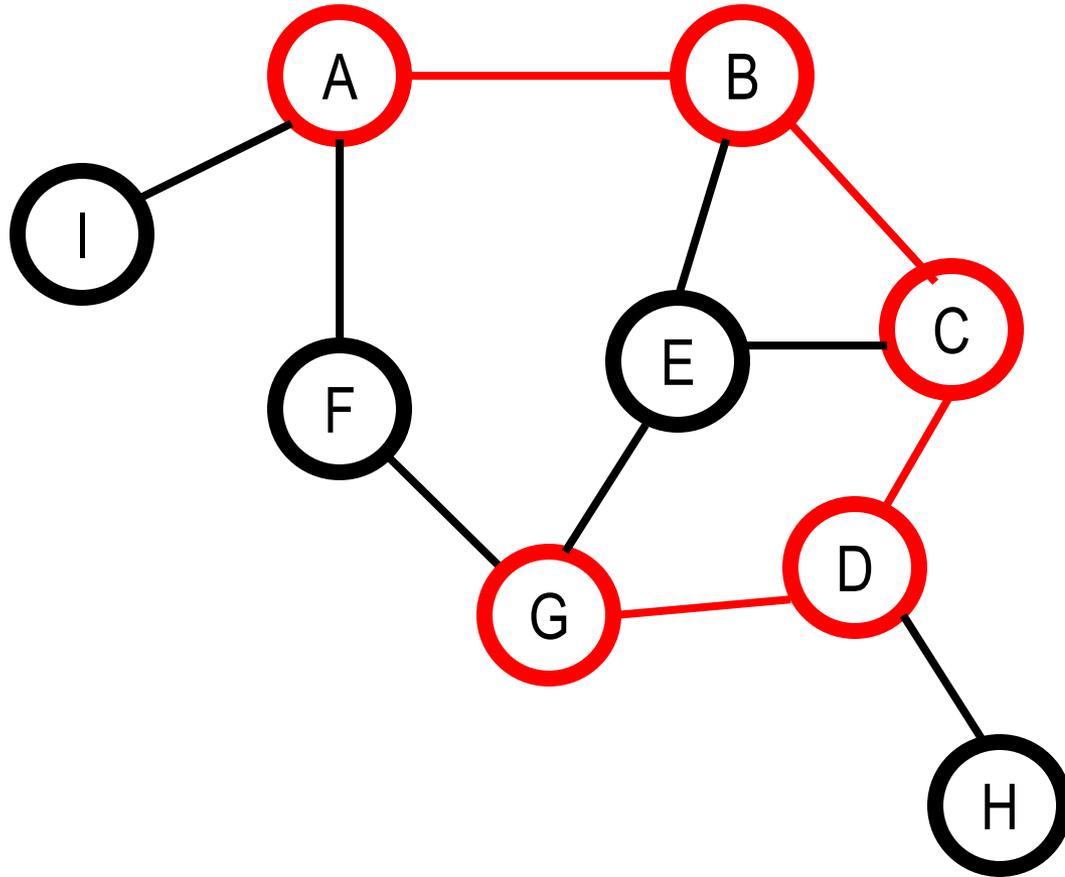
Example: DFS(A)



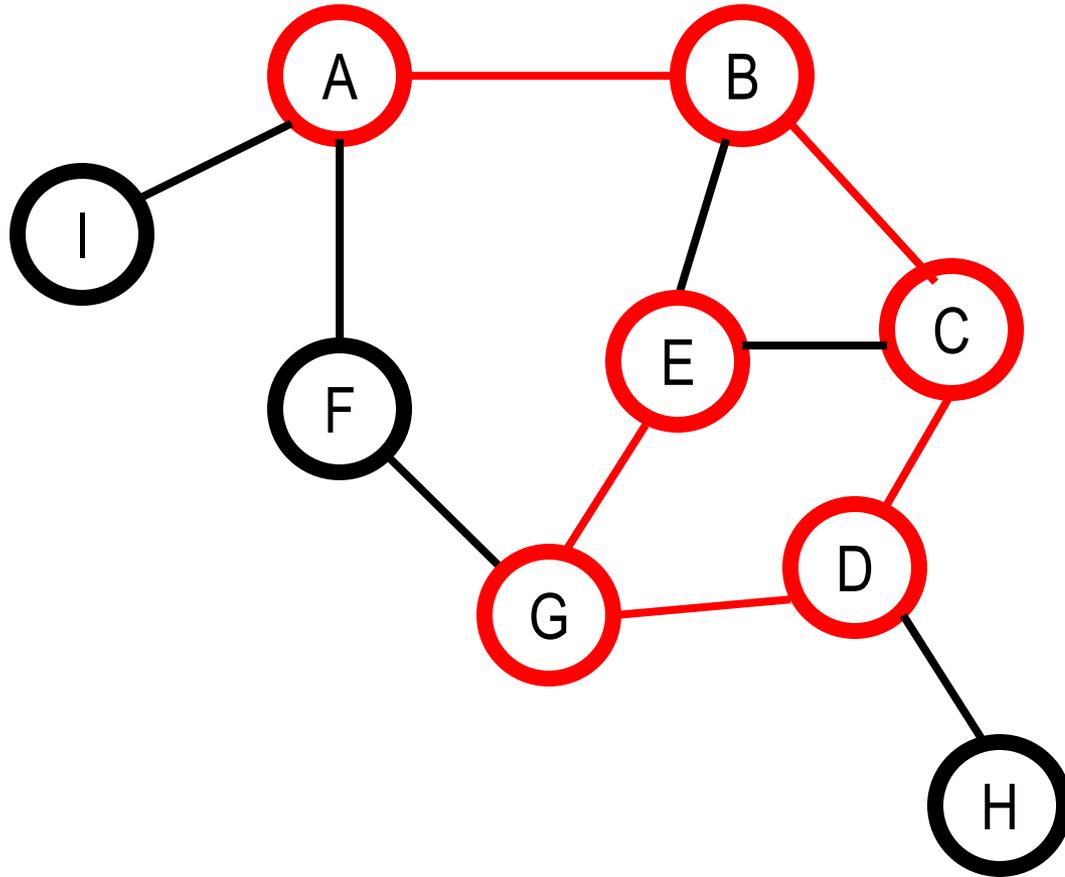
Example: DFS(A)



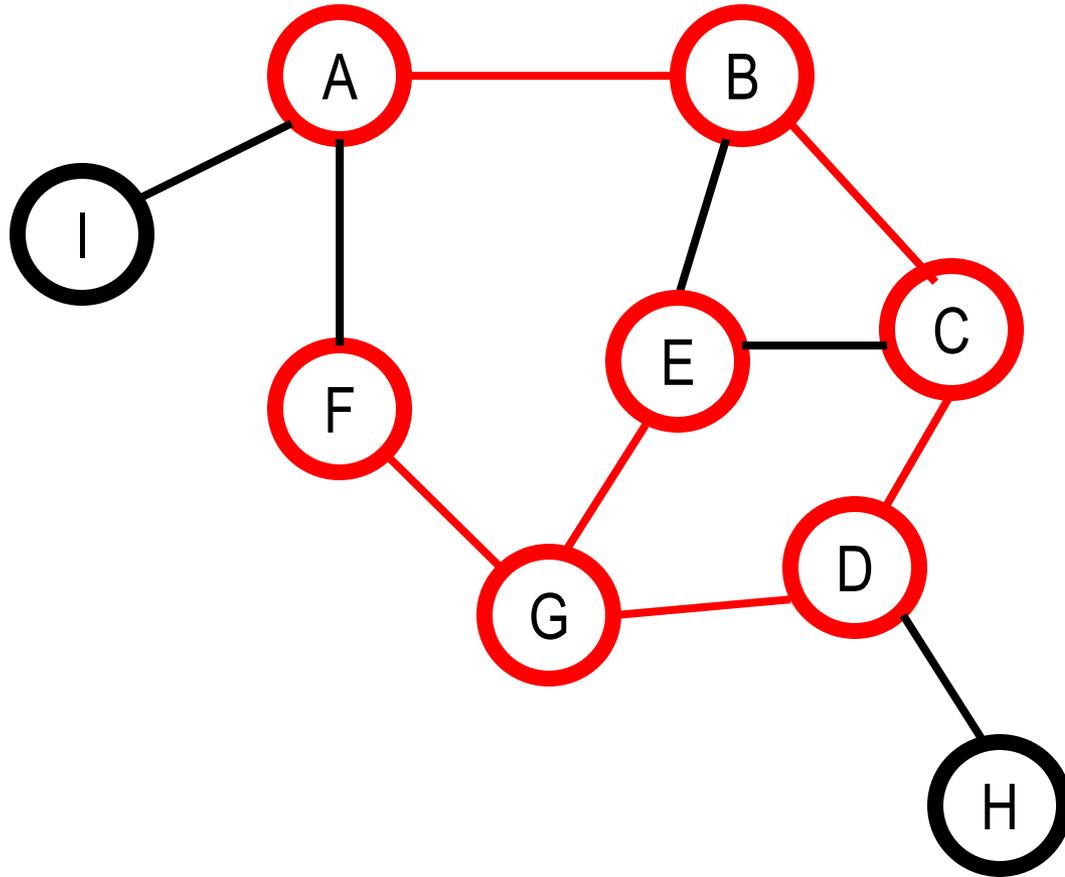
Example: DFS(A)



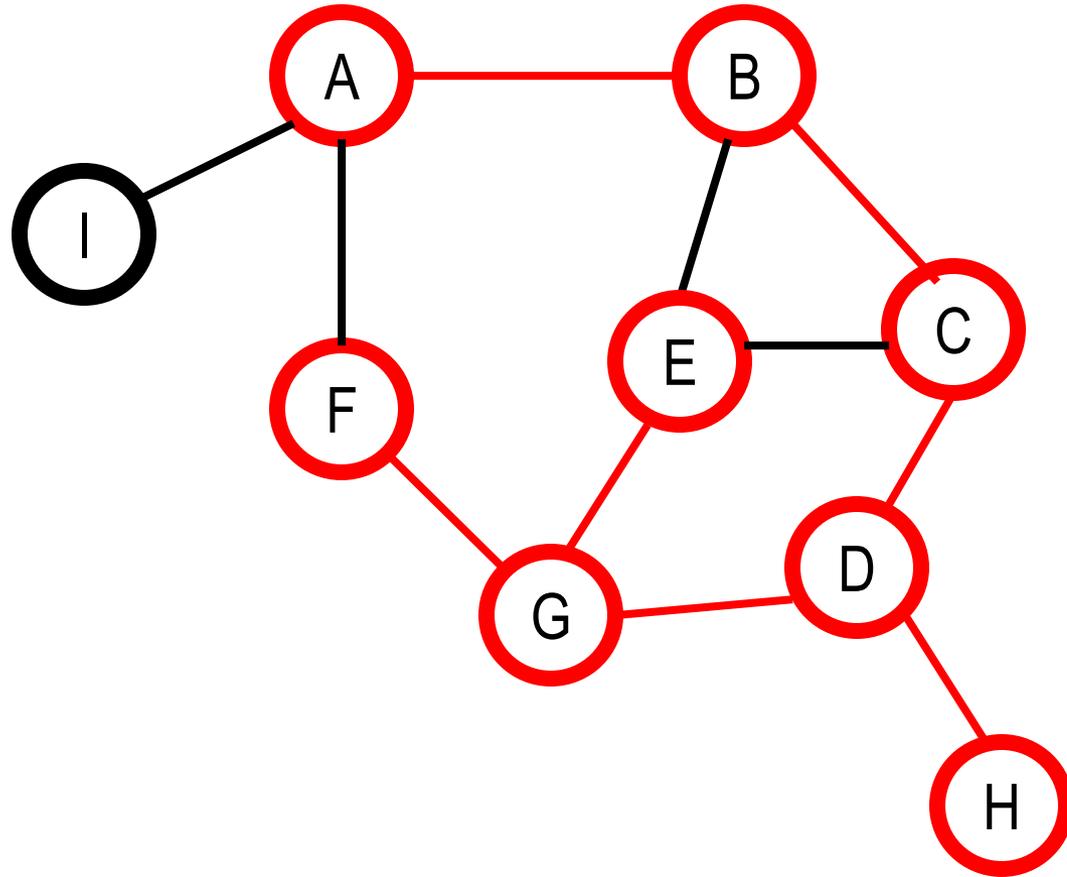
Example: DFS(A)



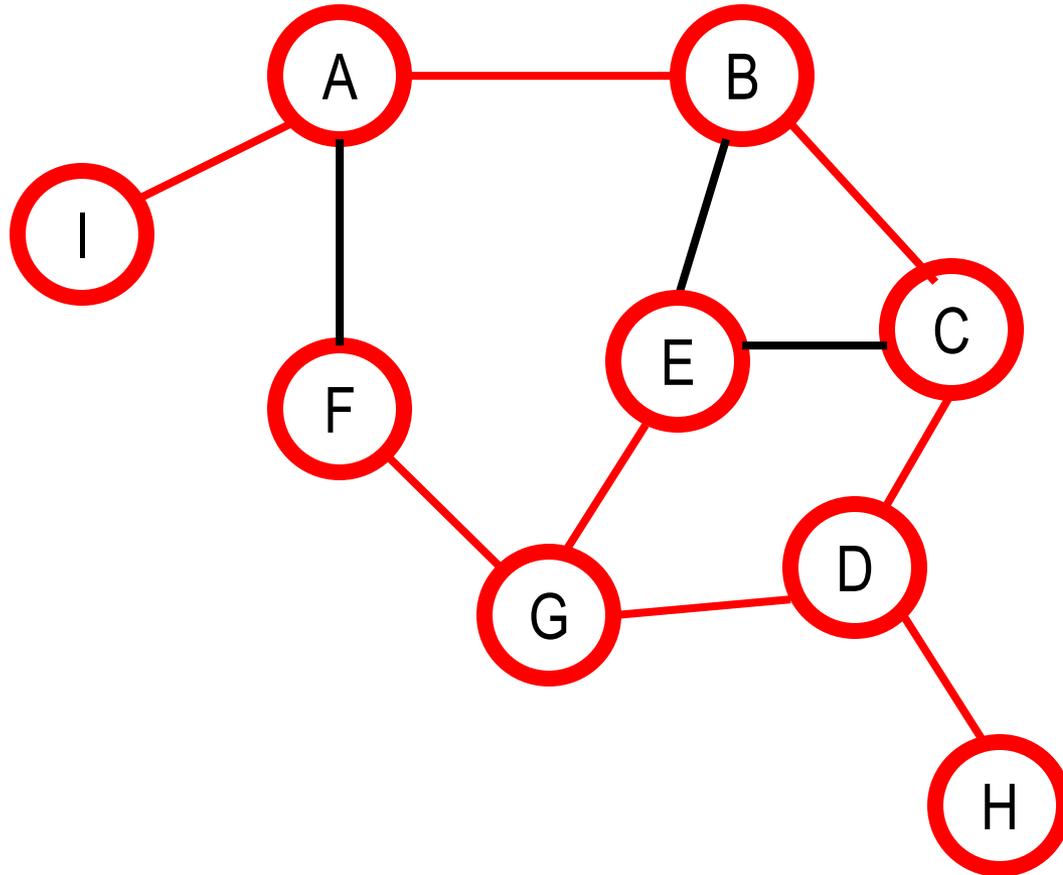
Example: DFS(A)



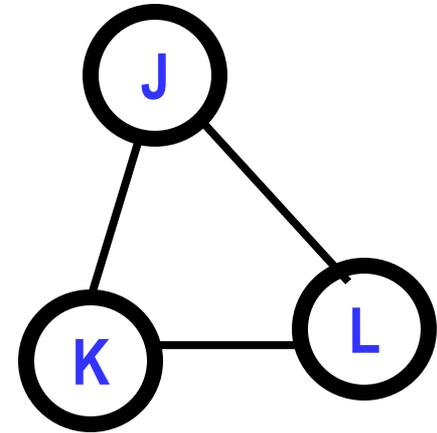
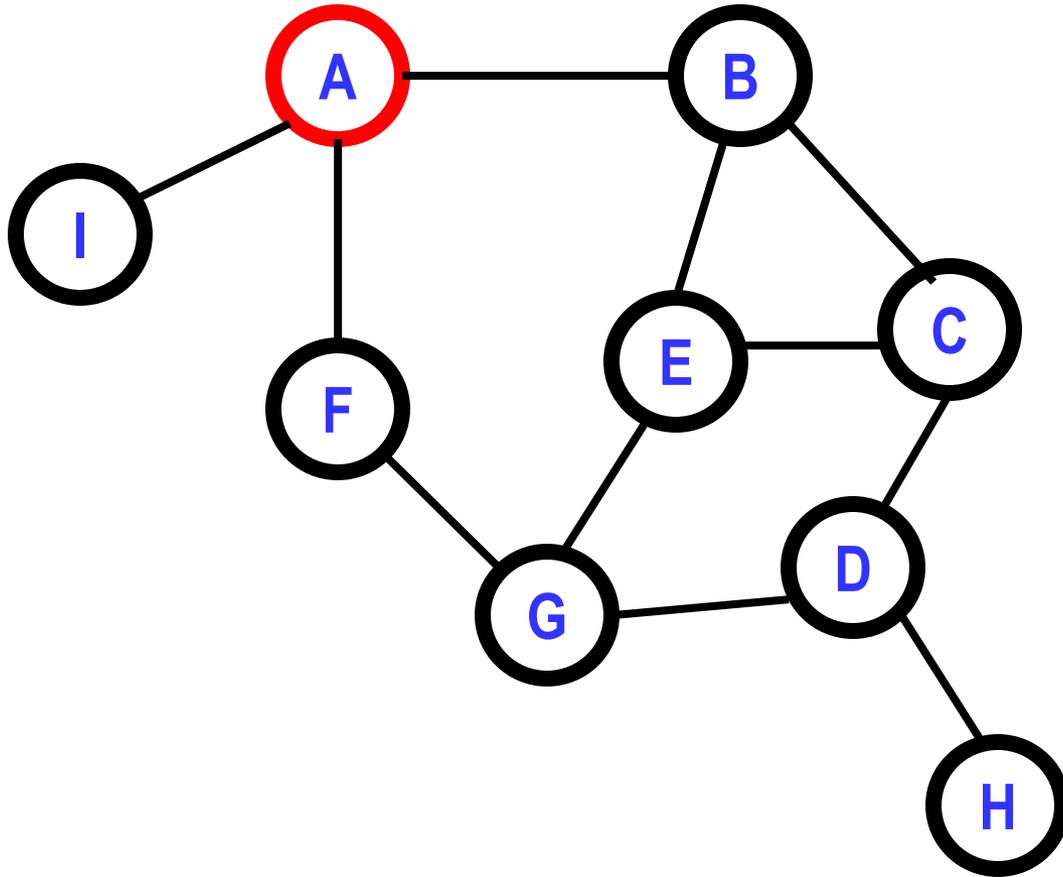
Example: DFS(A)



Example: DFS(A)



▶ QUIZ?



Depth-First Search (DFS) - Analysis

- Worst-case running time
 - Using adjacent matrix
 - ☞ $O(|V|^2)$ time
 - Using adjacent list
 - ☞ $O(|V|+|E|)$ time

Depth-First Search (DFS) Visualization

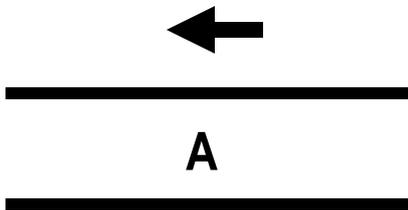
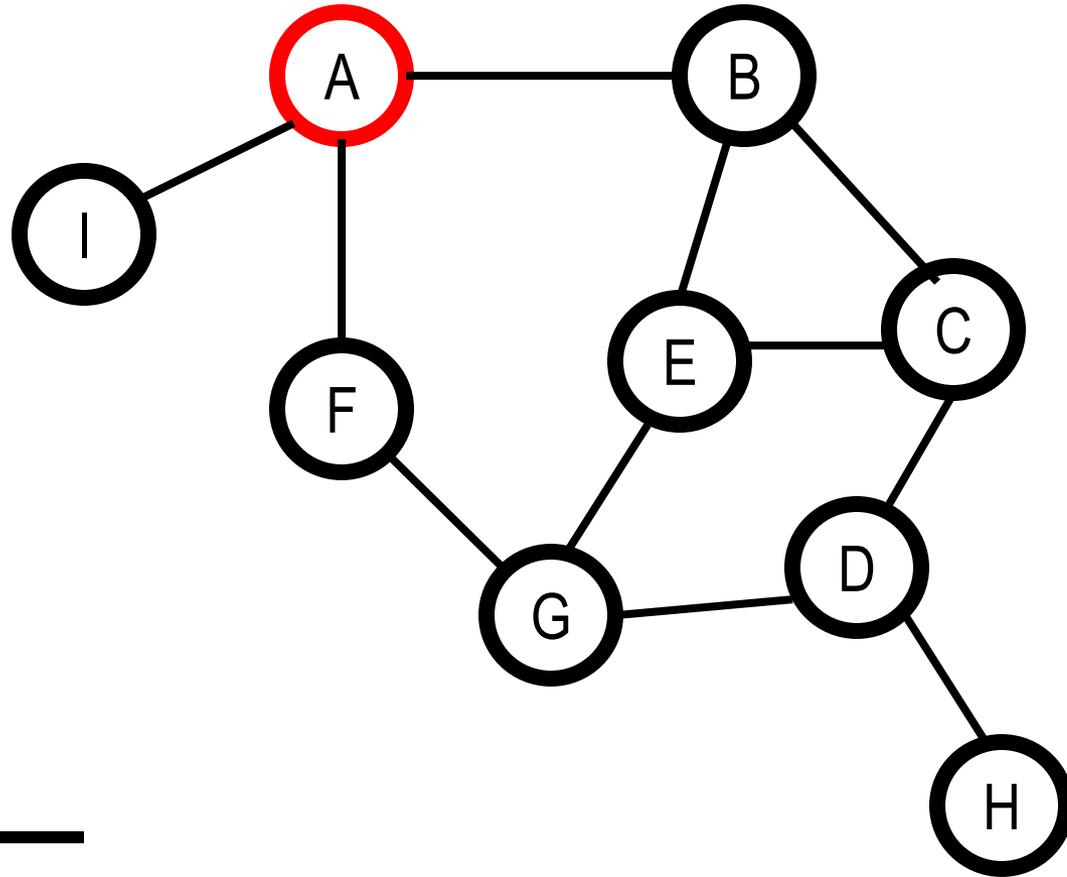
- *Depth-First Search (DFS) Visualization*



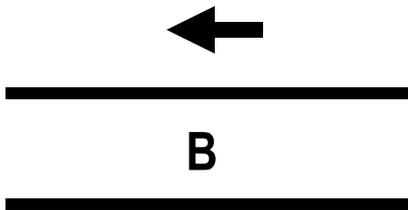
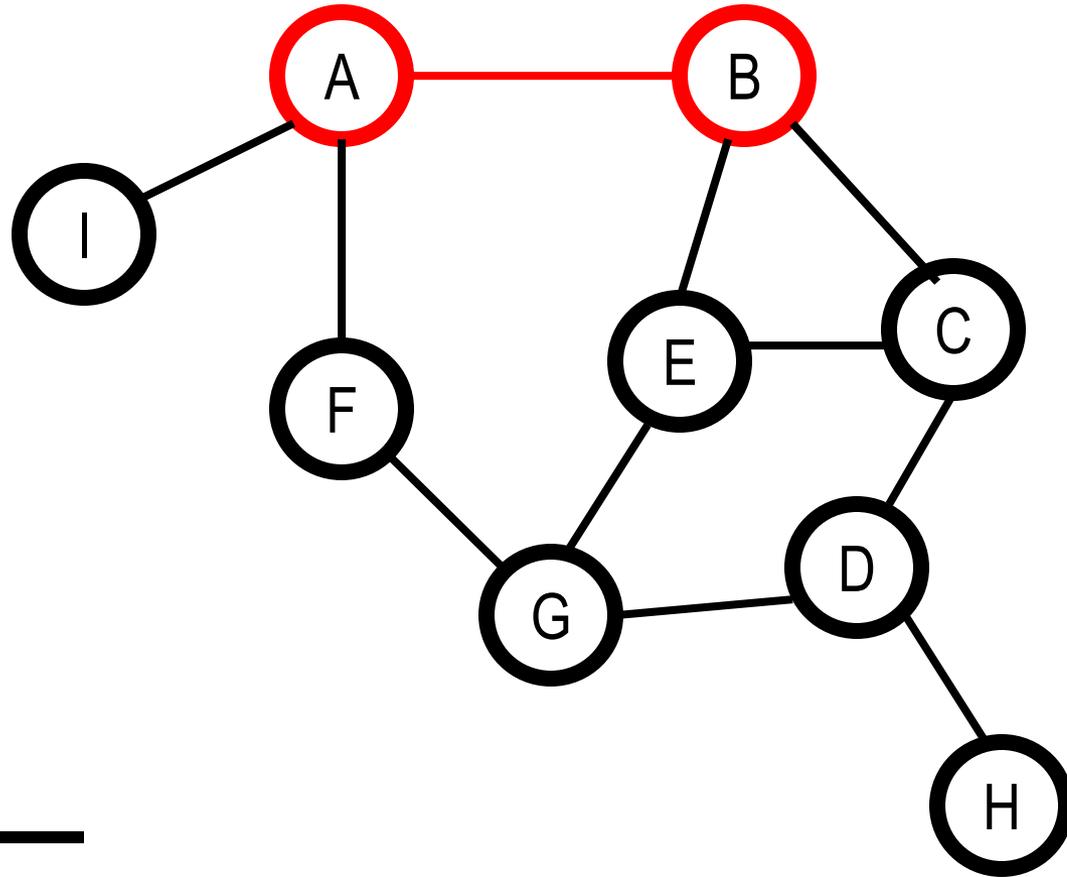
Breadth-First Search (BFS) – Iterative using Queue

- BFS(u):
 - Q.CreateQueue();
 - Q.Insert(u); Mark u as visited;
 - while !Q.QueueEmpty() do
 - ☞ Q.QueueDelete(w);
 - ☞ for each **unvisited** vertex v adjacent to w do
 - Mark v as visited;
 - Q.QueueInsert(v);

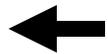
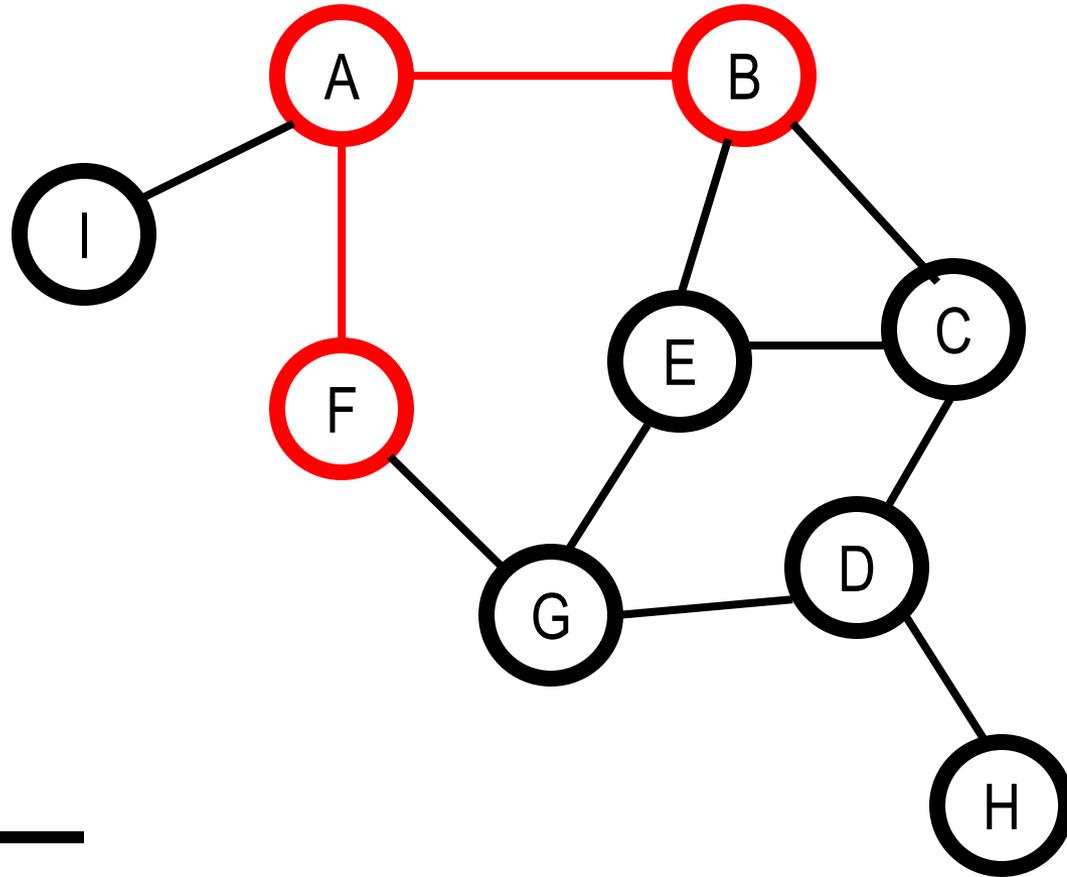
Example: BFS(A) - Iterative



Example: BFS(A) - Iterative

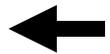
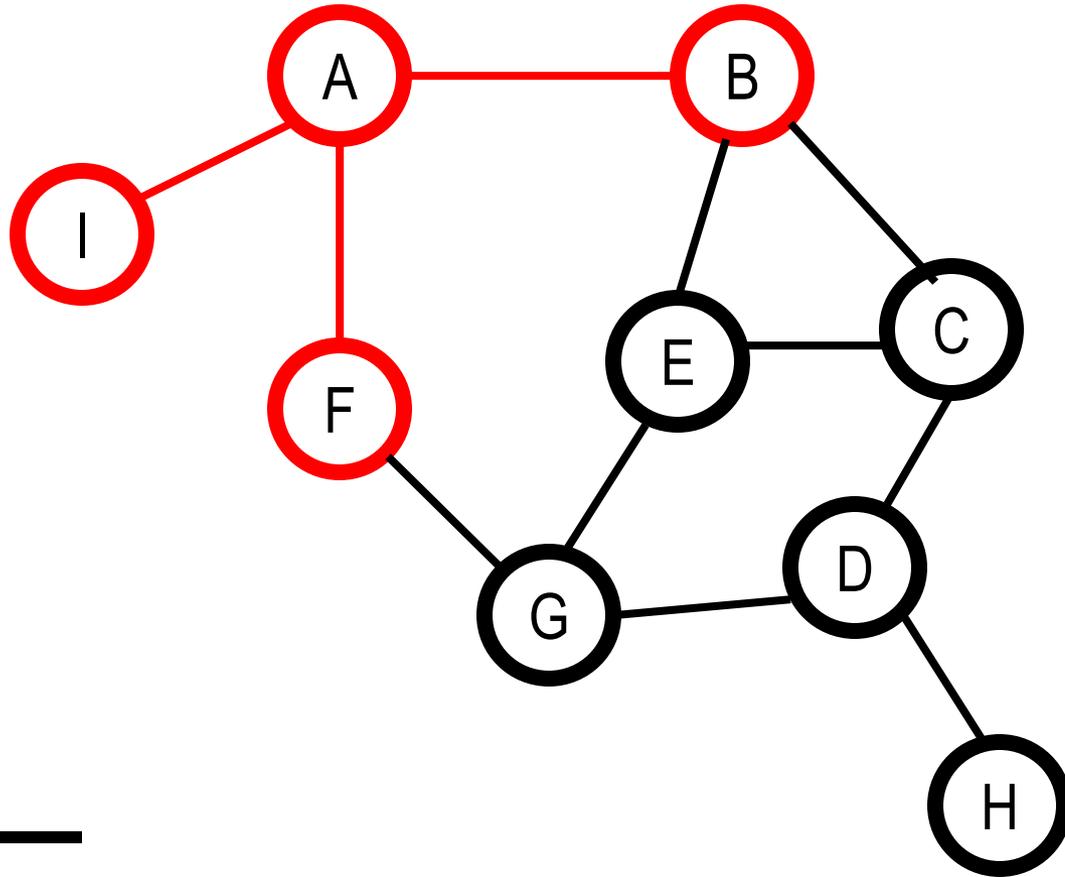


Example: BFS(A) - Iterative



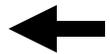
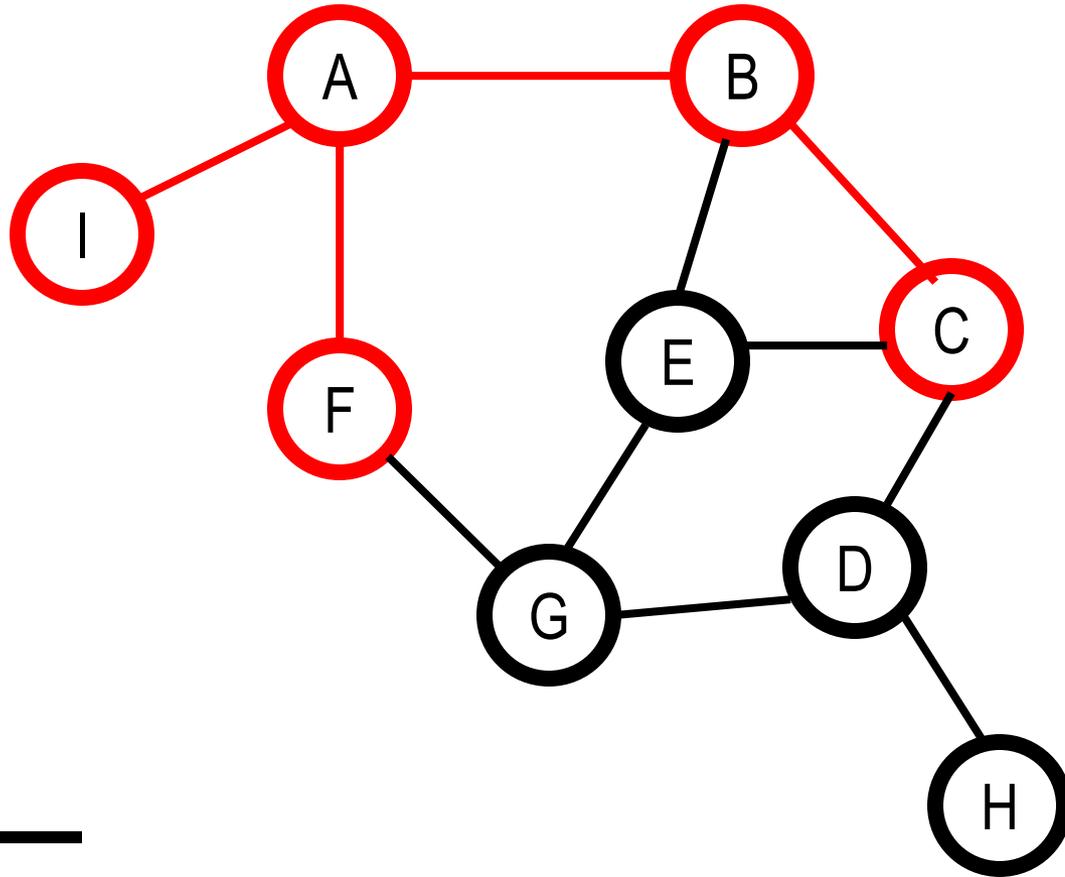
B F

Example: BFS(A) - Iterative



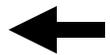
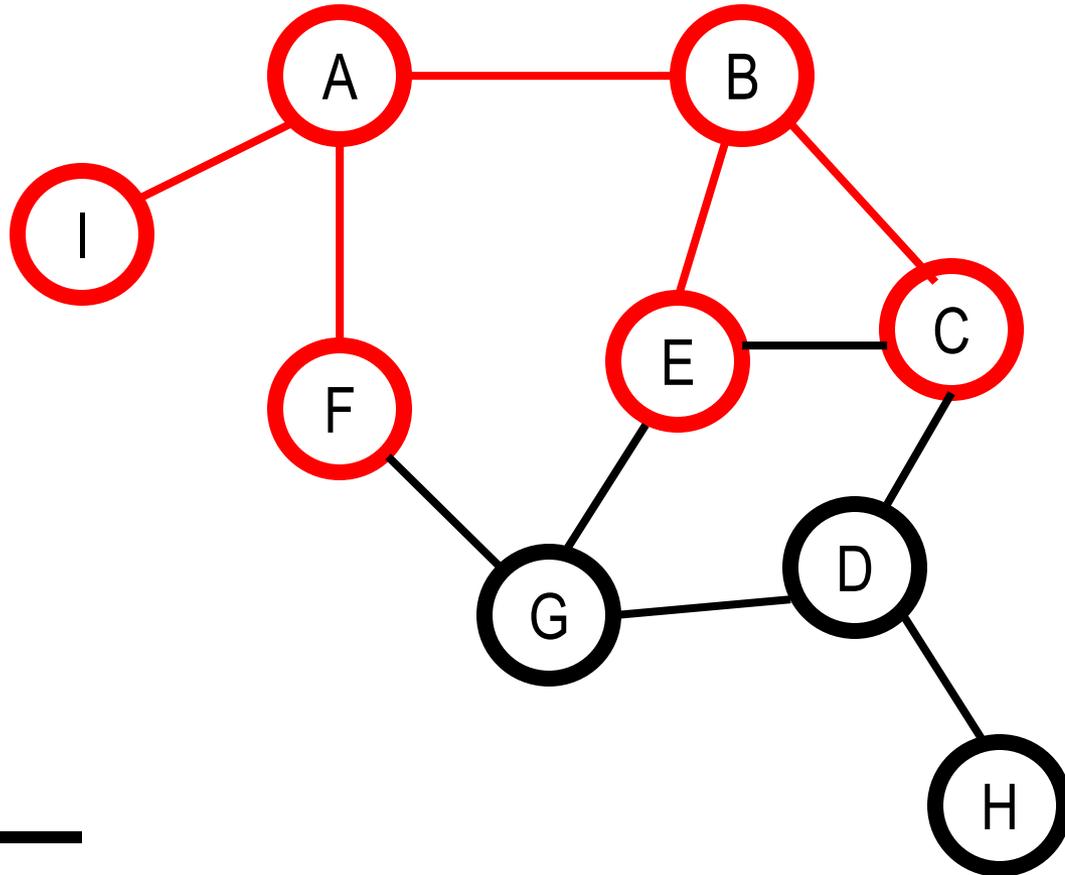
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Example: BFS(A) - Iterative



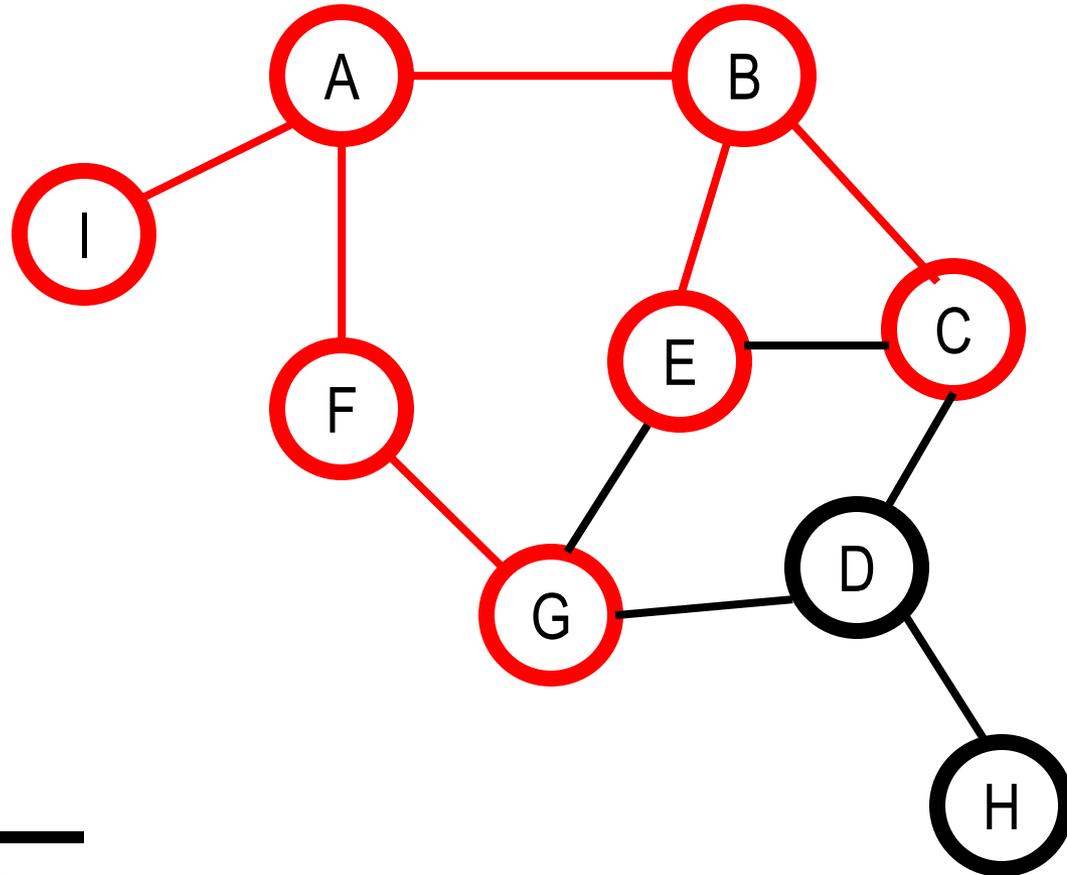
FIC

Example: BFS(A) - Iterative



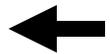
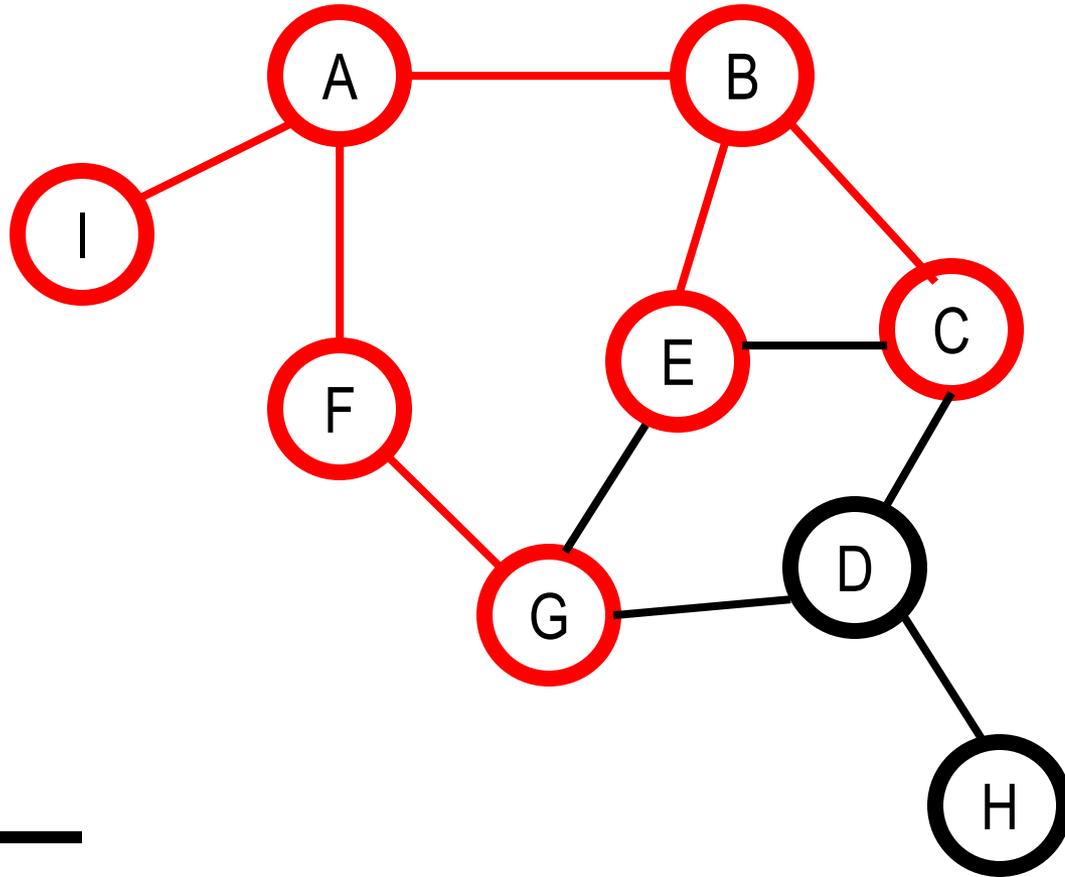
F I C E

Example: BFS(A) - Iterative



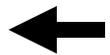
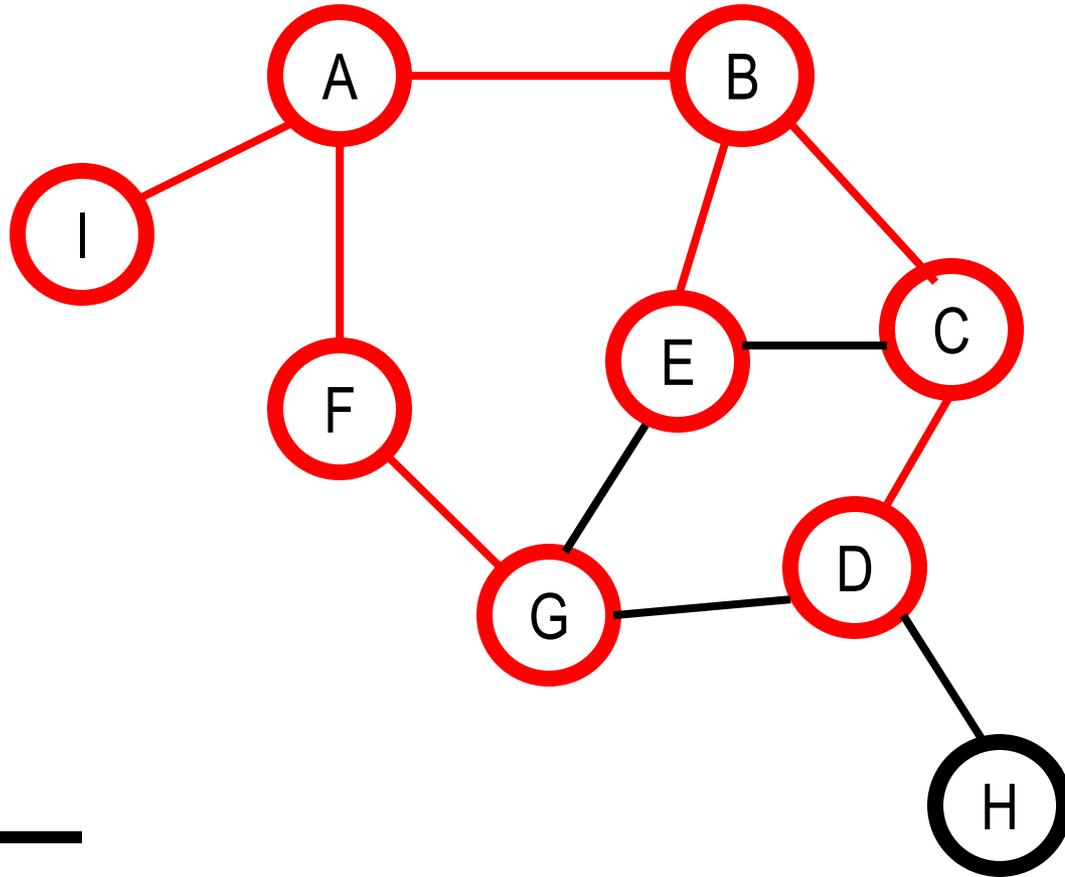
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Example: BFS(A) - Iterative



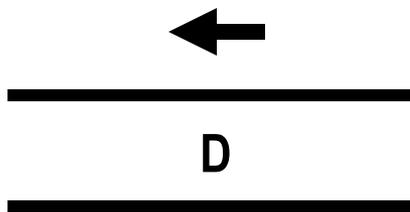
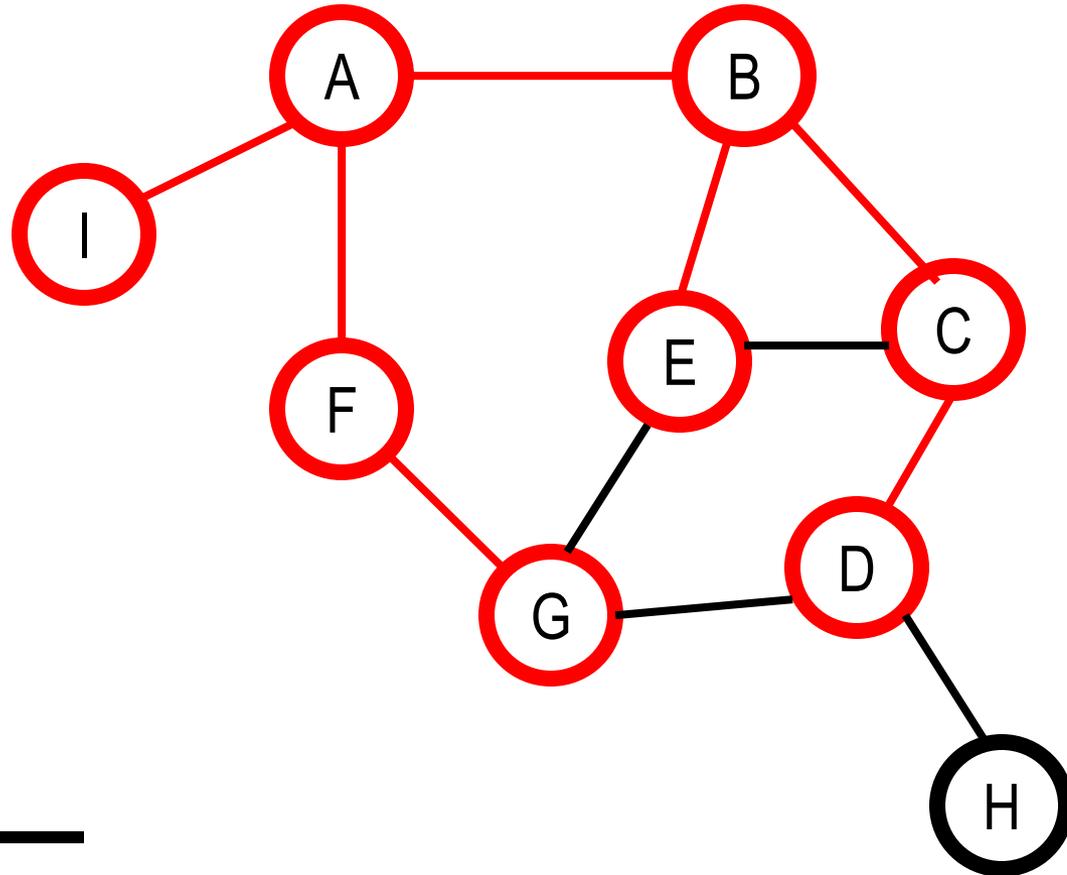
C E G

Example: BFS(A) - Iterative

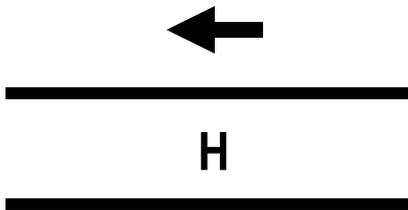
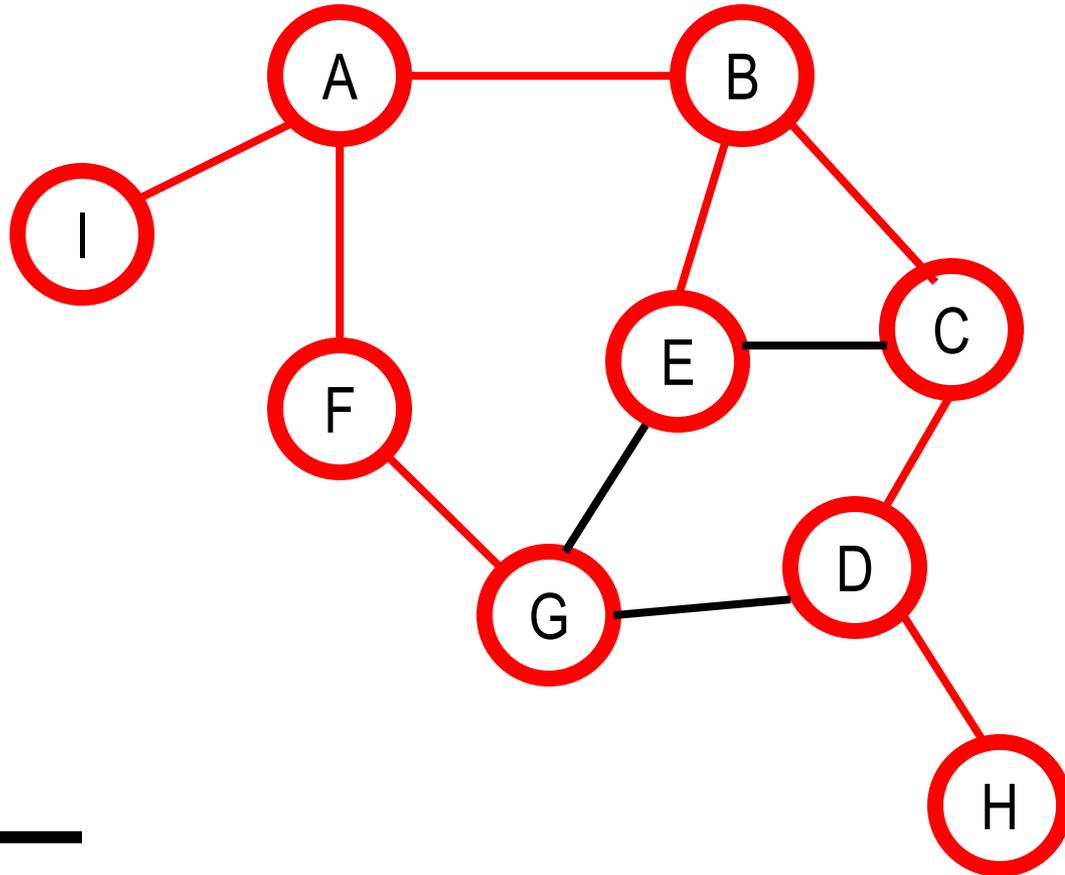


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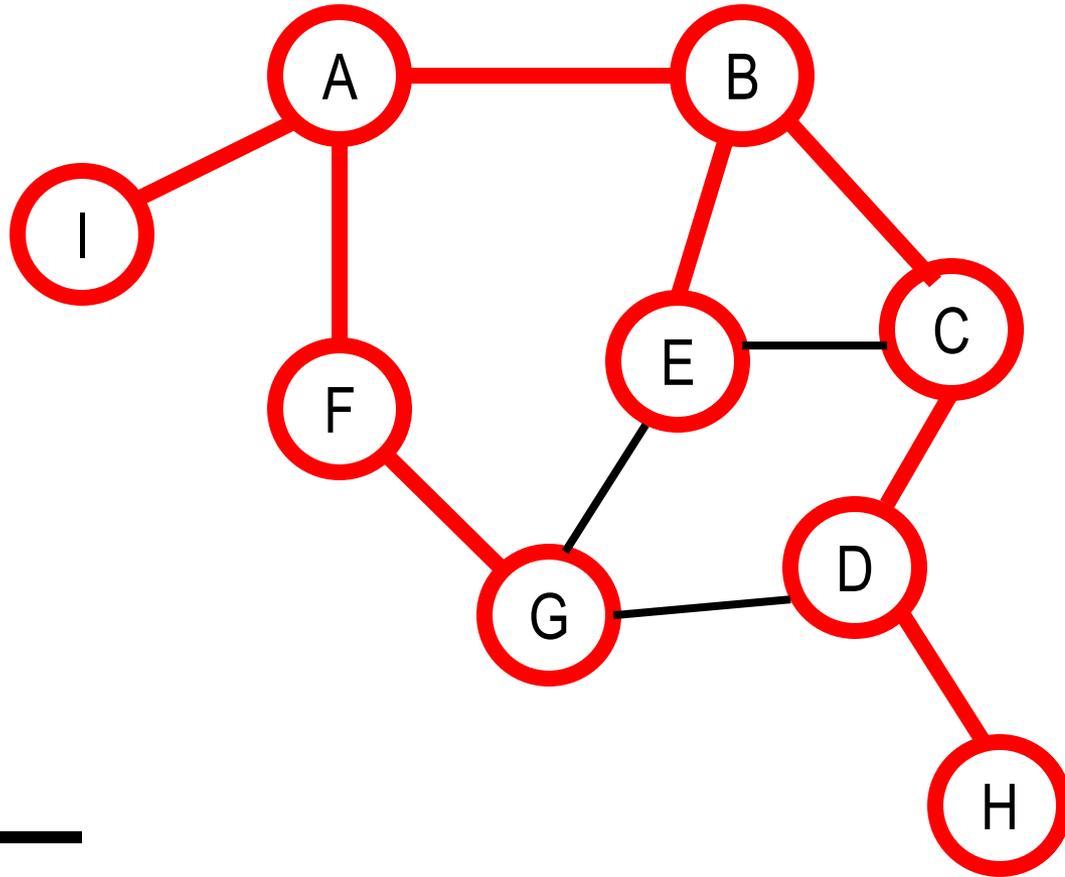
Example: BFS(A) - Iterative



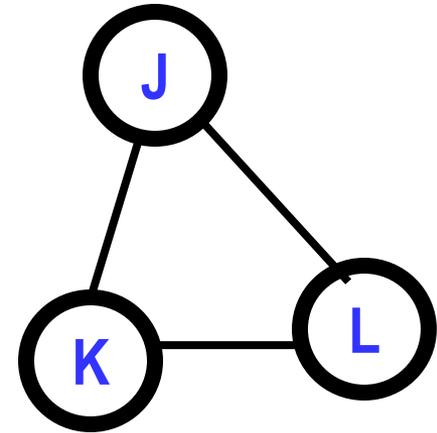
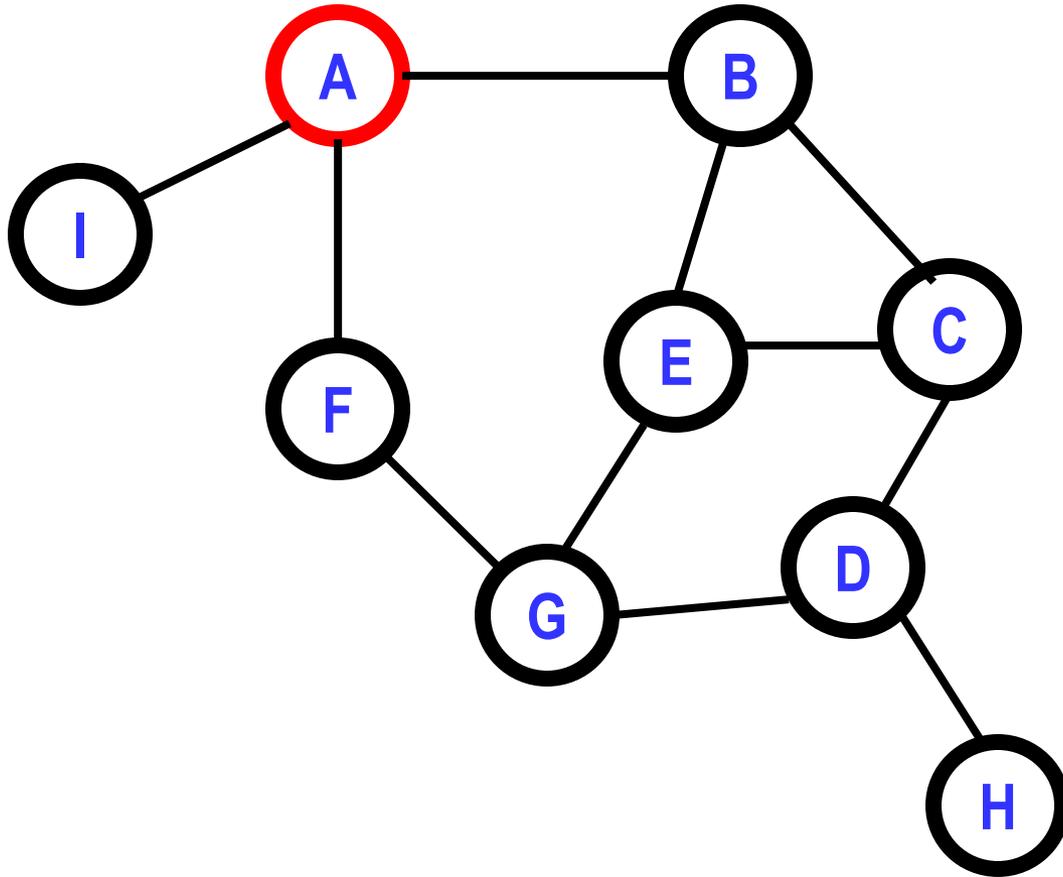
Example: BFS(A) - Iterative



Example: BFS(A) - Iterative



► QUIZ?



Breadth-First Search (BFS) - Analysis

- Worst-case running time
 - Using adjacent matrix
 - ☞ $O(|V|^2)$ time
 - Using adjacent list
 - ☞ $O(|V|+|E|)$ time

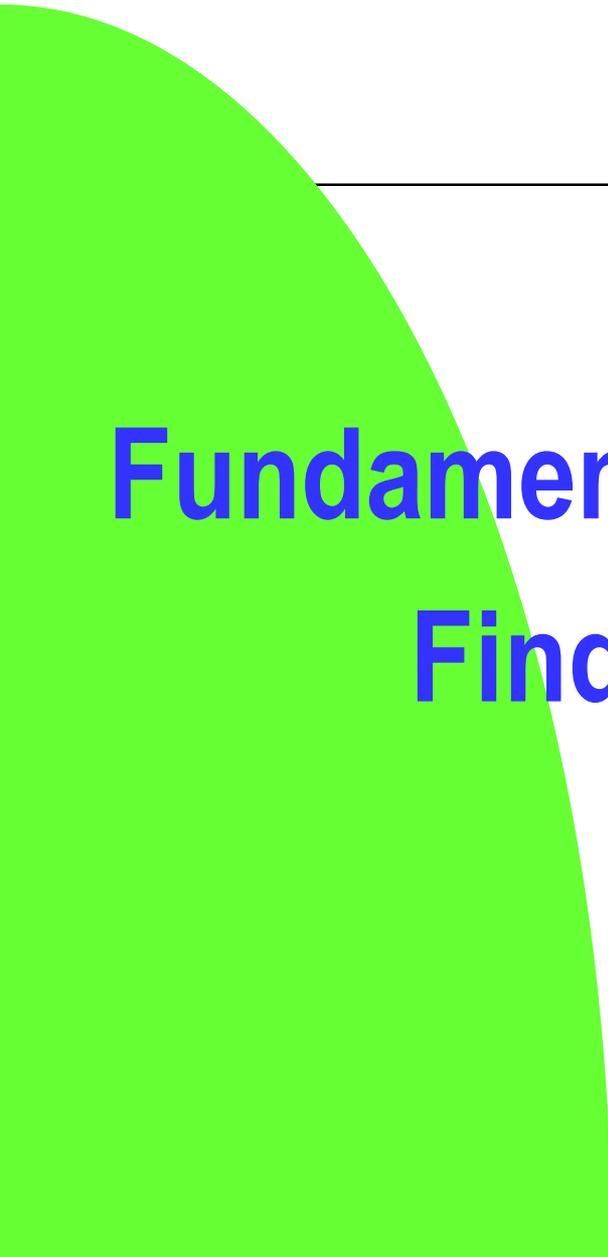
Breadth-First Search (BFS) Visualization

- *Breadth-First Search (BFS) Visualization*



► QUIZ?

- Compare graph traversal DFS vs. BFS?
- Which data structures are used for DFS and BFS of a graph?
- Best-First Search? (**Branch and Bound**)



Fundamental Problems on Graphs: Finding Shortest Paths

A Shortest-path From U to V in a Weighted, Directed Graph

- Let $G=(V,E)$ be a **edge-weighted, directed** graph with weight function $W: E \rightarrow \mathbb{R}$ mapping edges to real valued weights.
 - The length of a path from u to v is **the number of its constituent edges**.
 - The weighted length of a path P from u to v is the sum of the **weights** of its constituent edges.

A Shortest-path From U to V in a Weighted, Directed Graph

- The **shortest-path** weighted length from u to v is
 - Min (weights of all paths from u to v) if there is a path from u to v
 - ∞ otherwise
- A shortest path from u to v is **any path with the shortest-path weighted length.**

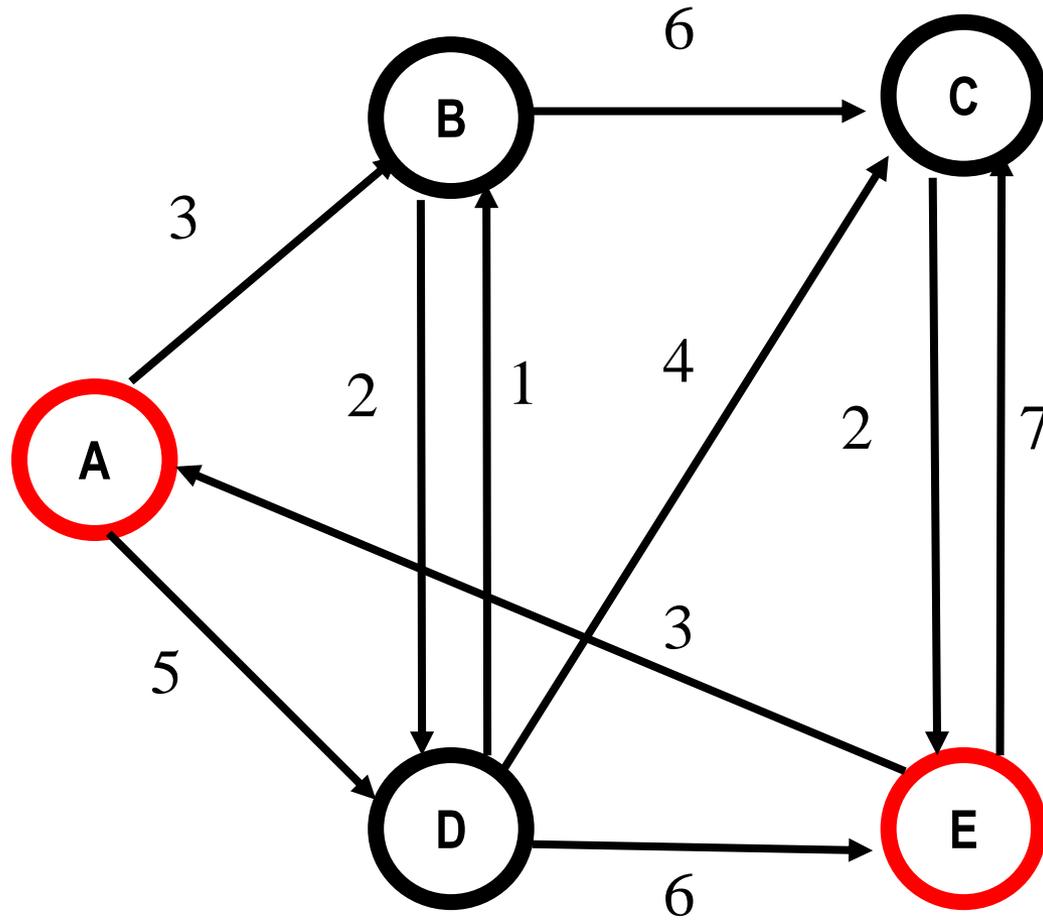
✓ Shortest Path Problems

- **The Single-Pair Shortest Path Problem**
- **The Single-Source Shortest Paths Problem**
- **The Single-Destination Shortest Paths Problem**
- **The All-Pairs Shortest Paths Problem**

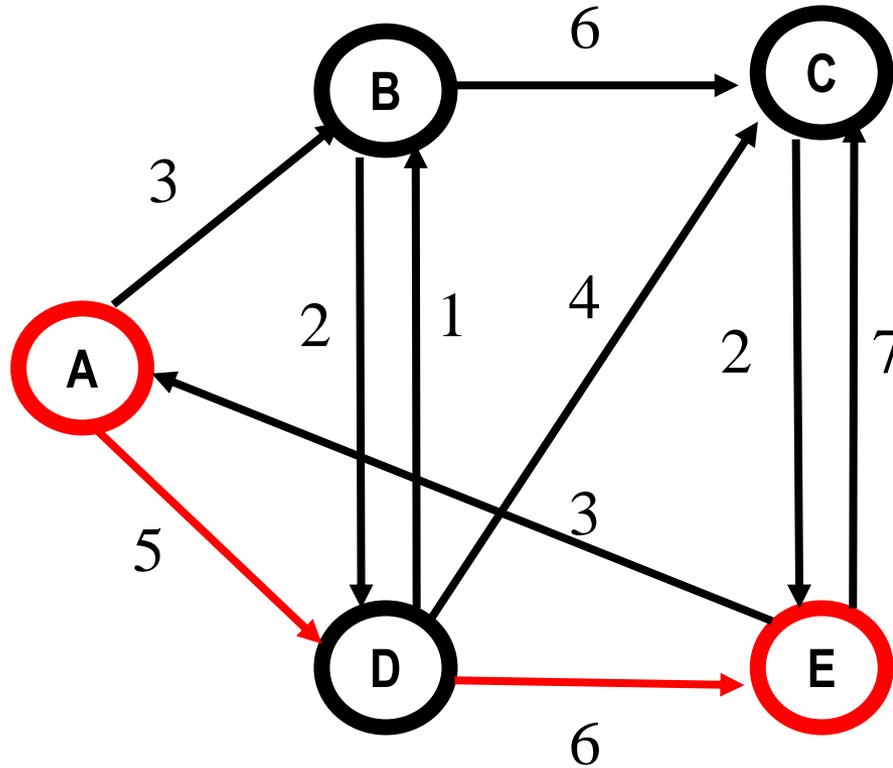
1. The Single-Pair Shortest Path Problem

- Given a weighted, directed graph $G=(V,E)$ with weight function $W: E \rightarrow \mathbb{R}$ mapping edges to real valued weights,
- Find a shortest path from a given source vertex s to a given destination vertex d in V .
 - The Single-Pair Shortest Path Problem

Example: Shortest-Paths from A to E?

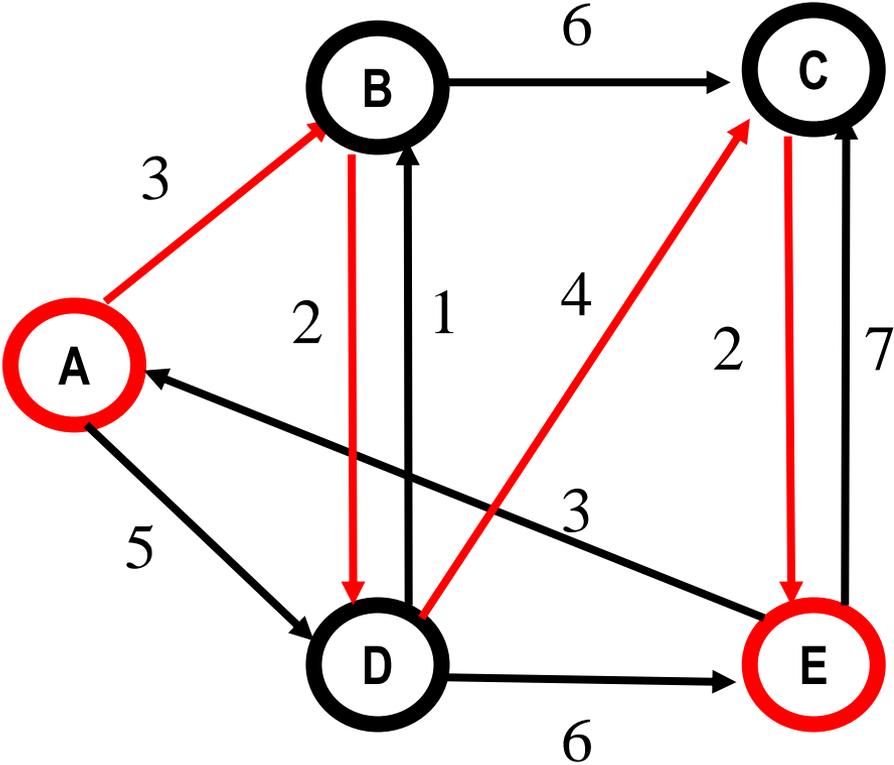


Example: A Shortest-Path from A to E



P1 = A, D, E 11

Example: Another Shortest-Path from A to E

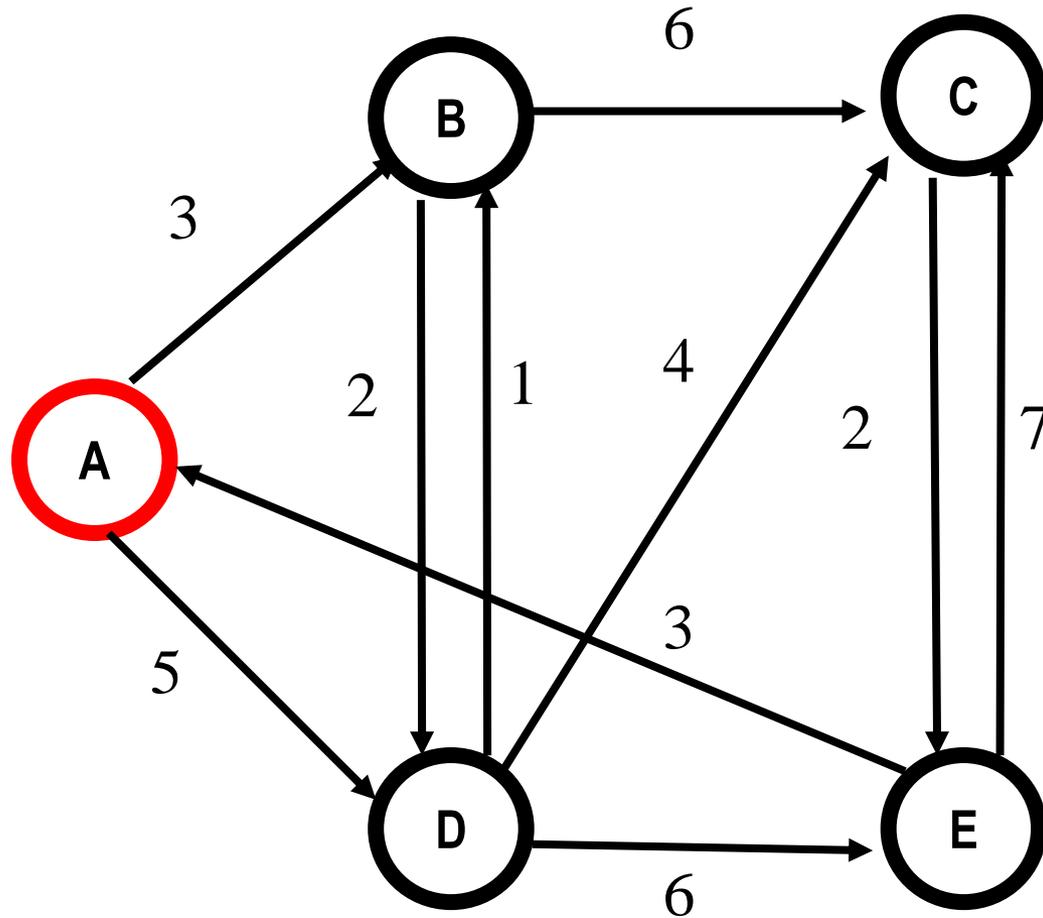


P2 = A, B, D, C, E 11

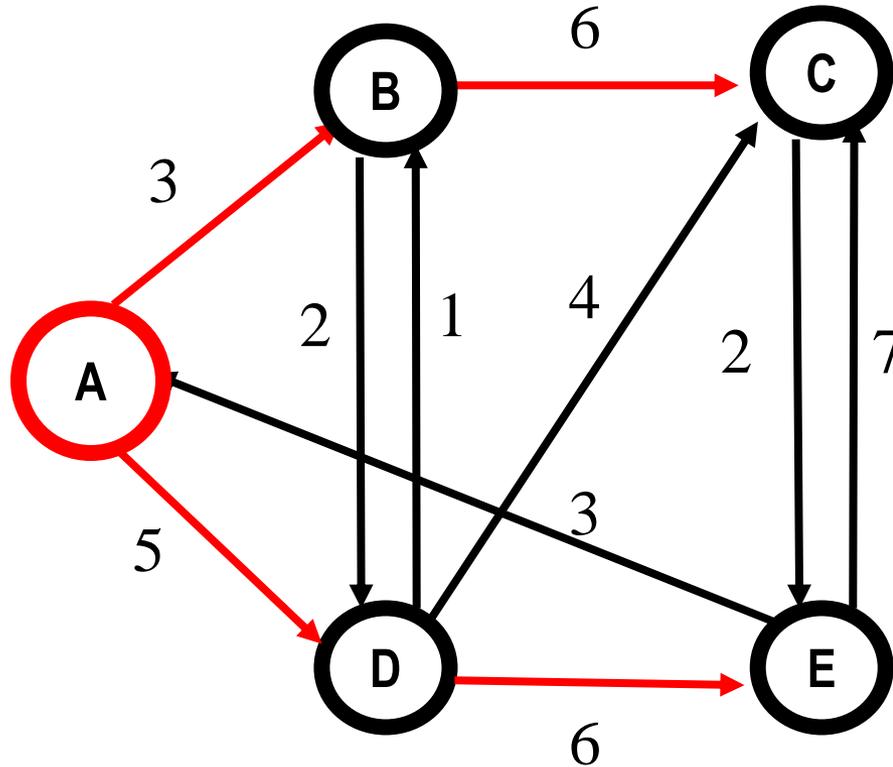
2. The Single-Source Shortest Paths Problem

- Given a weighted, directed graph $G=(V,E)$ with weight function $W: E \rightarrow \mathbb{R}$ mapping edges to real valued weights,
- Find a shortest path from a given source vertex s to every vertex v in V .
 - The Single-Source Shortest Paths Problem
 - The Single-Source/All-Destinations Shortest Paths Problem

Example: Shortest-Paths from A?

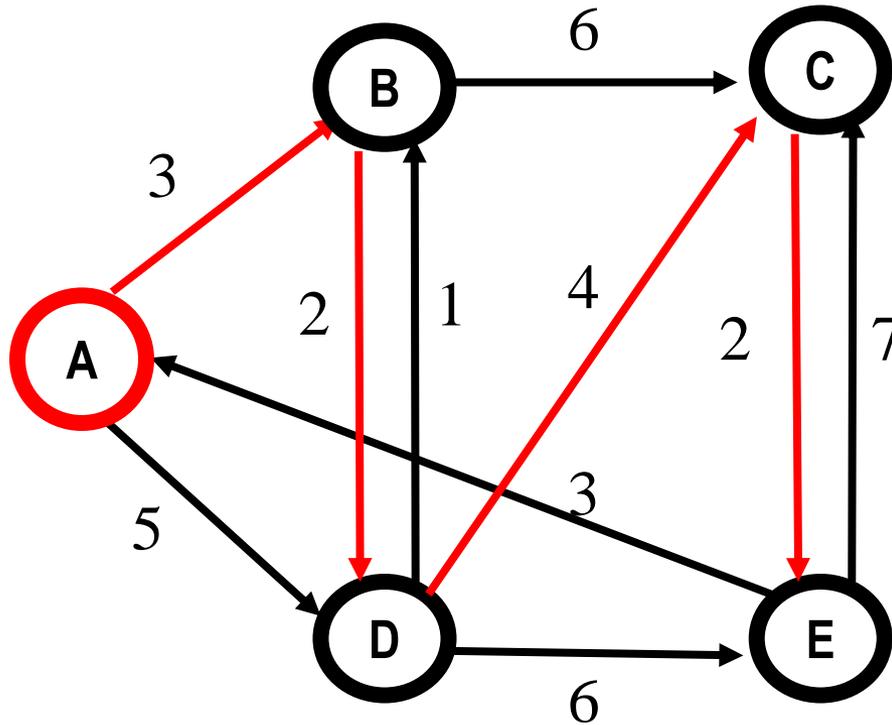


Example: Shortest-Path from A



A	A	B	C	D	E
	0	3	9	5	11

Example: Shortest-Path from A



A

A	B	C	D	E
0	3	9	5	11

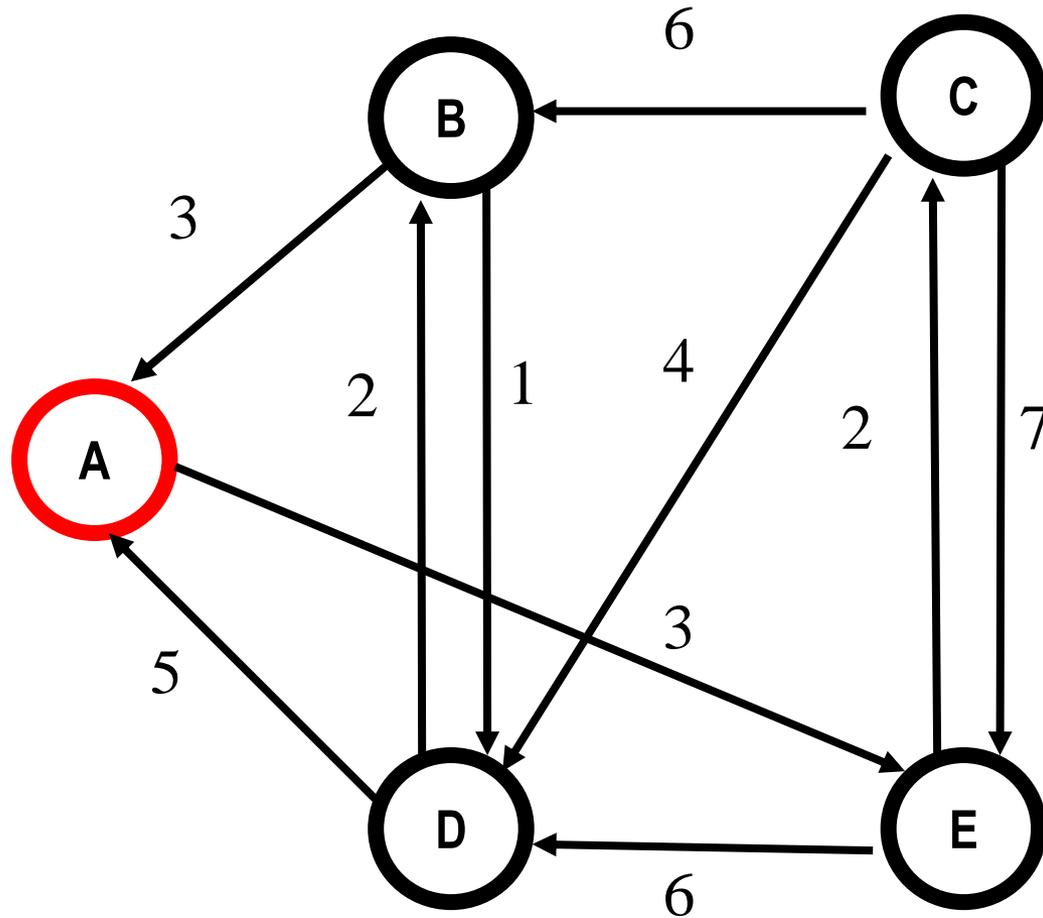
The Single-Source Shortest Paths Problem

- The **brute-force** algorithm is to consider **all** possible factorial paths and take the shortest.
- **Bellman-Ford Algorithm**
 - Dynamic Programming
- **Dijkstra Algorithm**
 - Greedy Approach

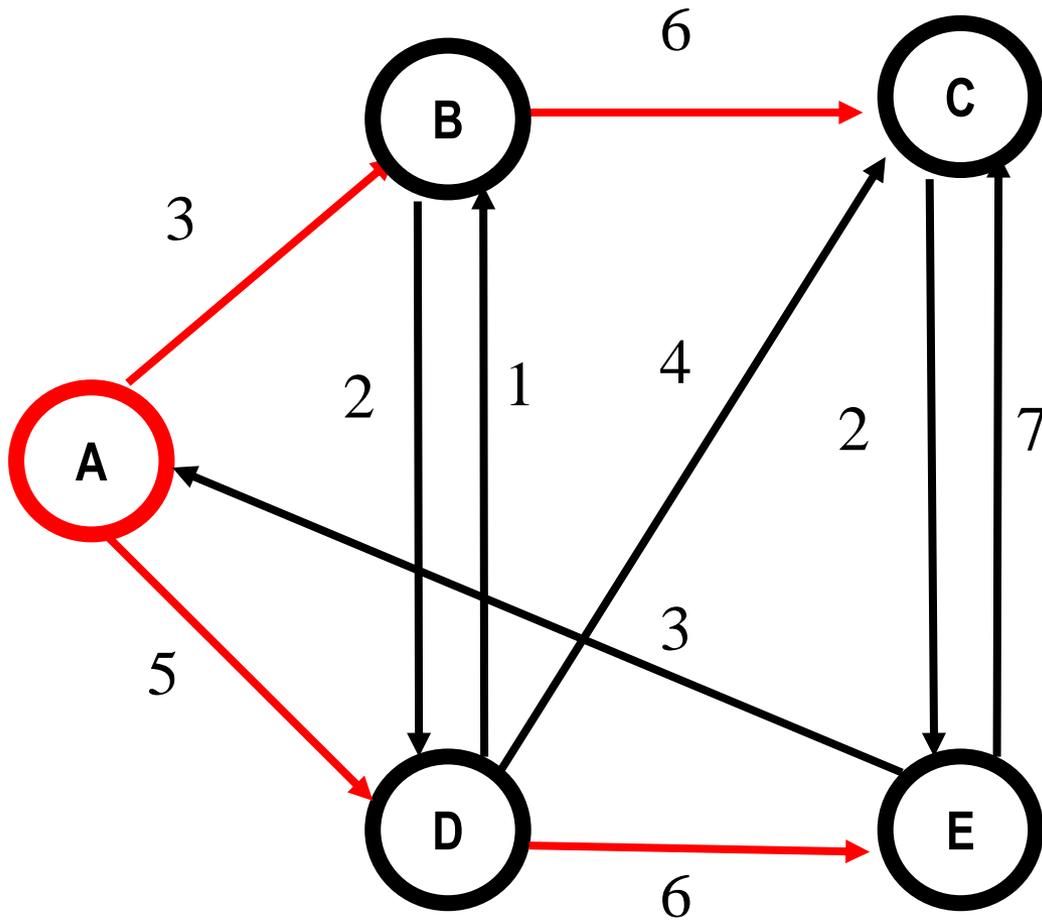
3. The Single-Destination Shortest Paths Problem

- Given a weighted, directed graph $G=(V,E)$ with weight function $W: E \rightarrow \mathbb{R}$ mapping edges to real valued weights,
- Find a shortest path to a given destination vertex d from every vertex v in V .
 - The Single-Destination Shortest Paths Problem
 - The Single-Destination/All-Sources Shortest Paths Problem
 - ☞ The reverse of the single-source shortest-paths problem!

Example: Shortest-Paths to A



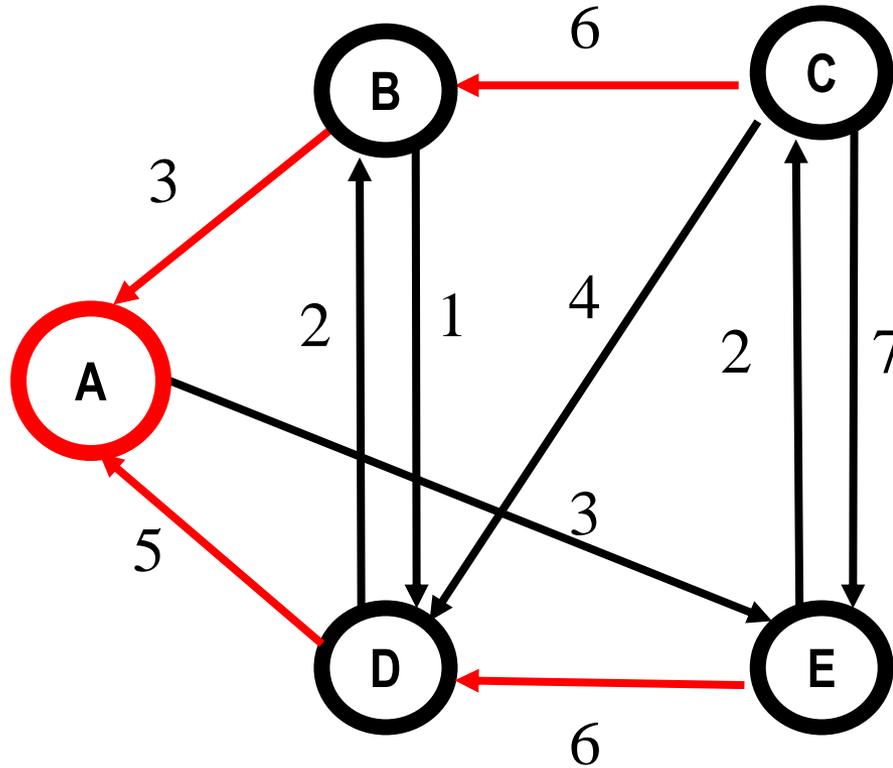
Example: Shortest-Paths from A



A

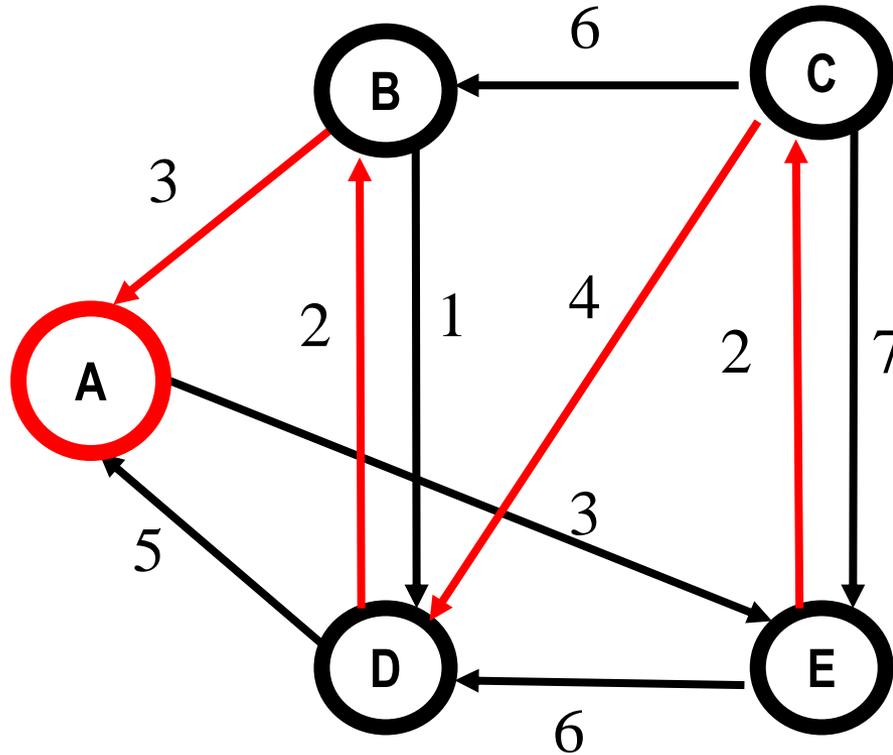
A	B	C	D	E
0	3	9	5	11

Example: Shortest-Paths to A



A	B	C	D	E
0	3	9	5	11

Example: Shortest-Paths to A

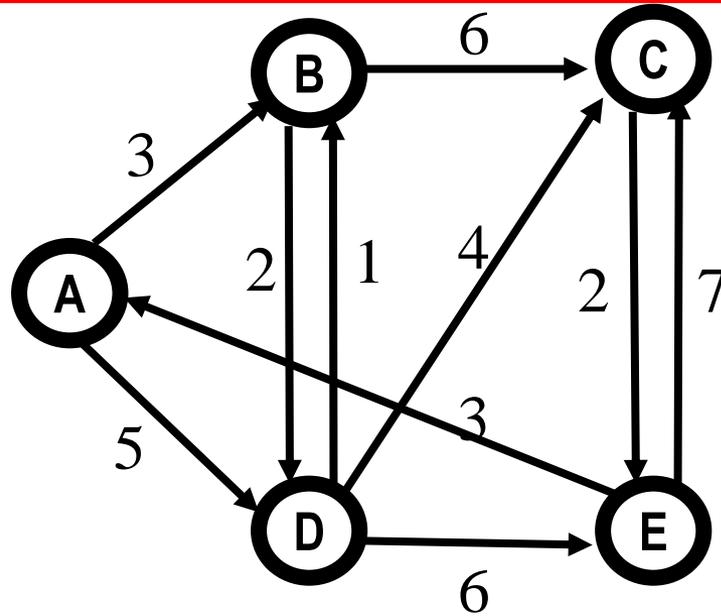


A	B	C	D	E
0	3	9	5	11

4. The All-Pairs Shortest Paths Problem

- Given a weighted, directed graph $G=(V,E)$ with weight function $W: E \rightarrow \mathbb{R}$ mapping edges to real valued weights,
- Find a shortest path from u to v for every pair of vertices u and v in V .
 - The All-Pairs Shortest Paths Problem

Example: Shortest-paths for Every Pair of Vertices



	A	B	C	D	E
A	0	3	9	5	11
B					
C					
D					
E					

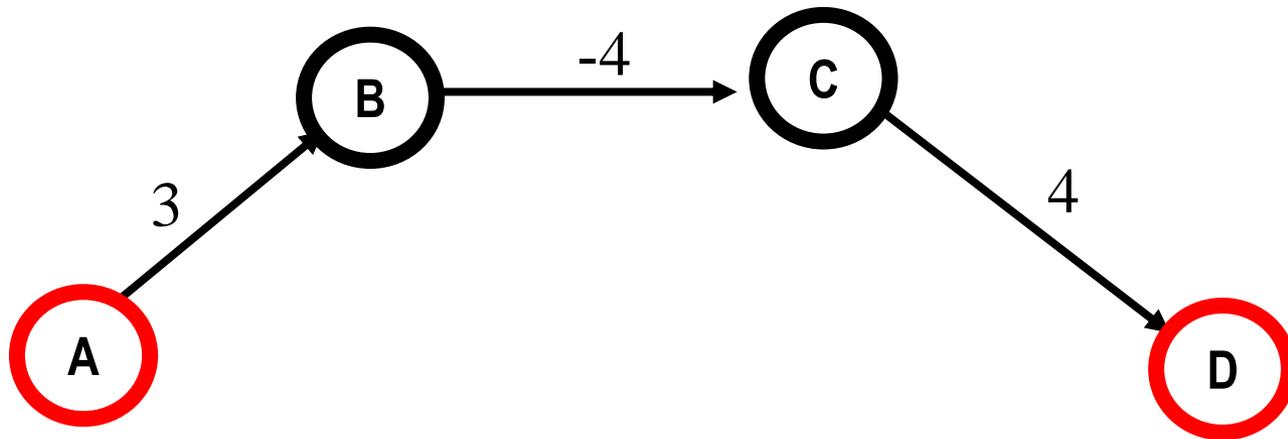
The All-Pairs Shortest Paths Problem

- The **brute-force** algorithm is to consider **all** possible factorial paths and take the shortest.
- **Floyd-Warshall Algorithm**
 - **Dynamic Programming**

Negative-Weight Edges? In The Shortest-Paths Problems

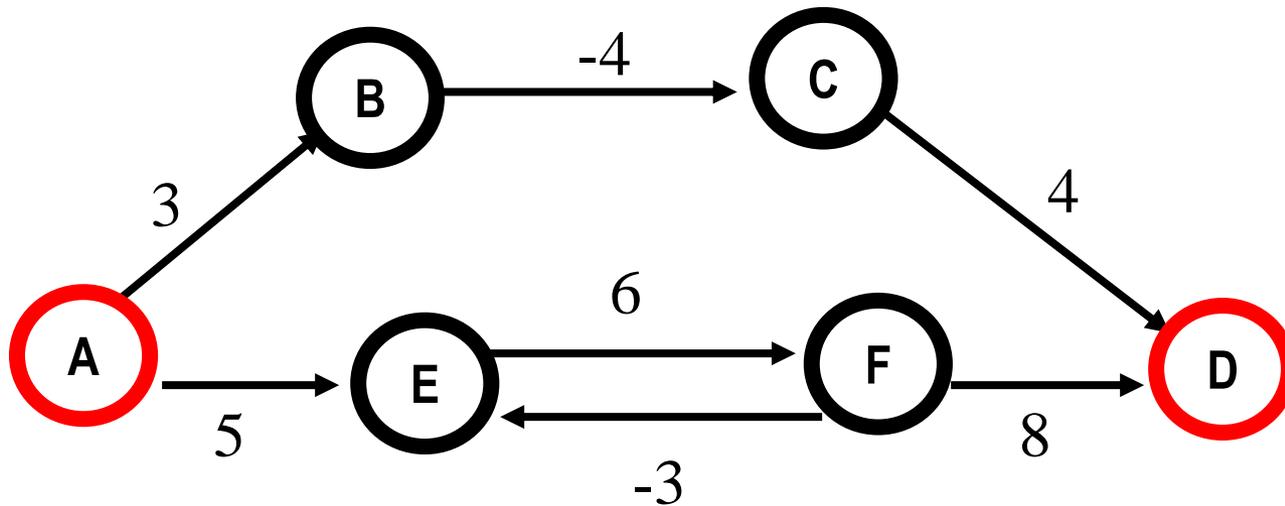
- Any problem with edges with **negative weights**?
 - Yes and No!

Example: Negative-Weight Edges?



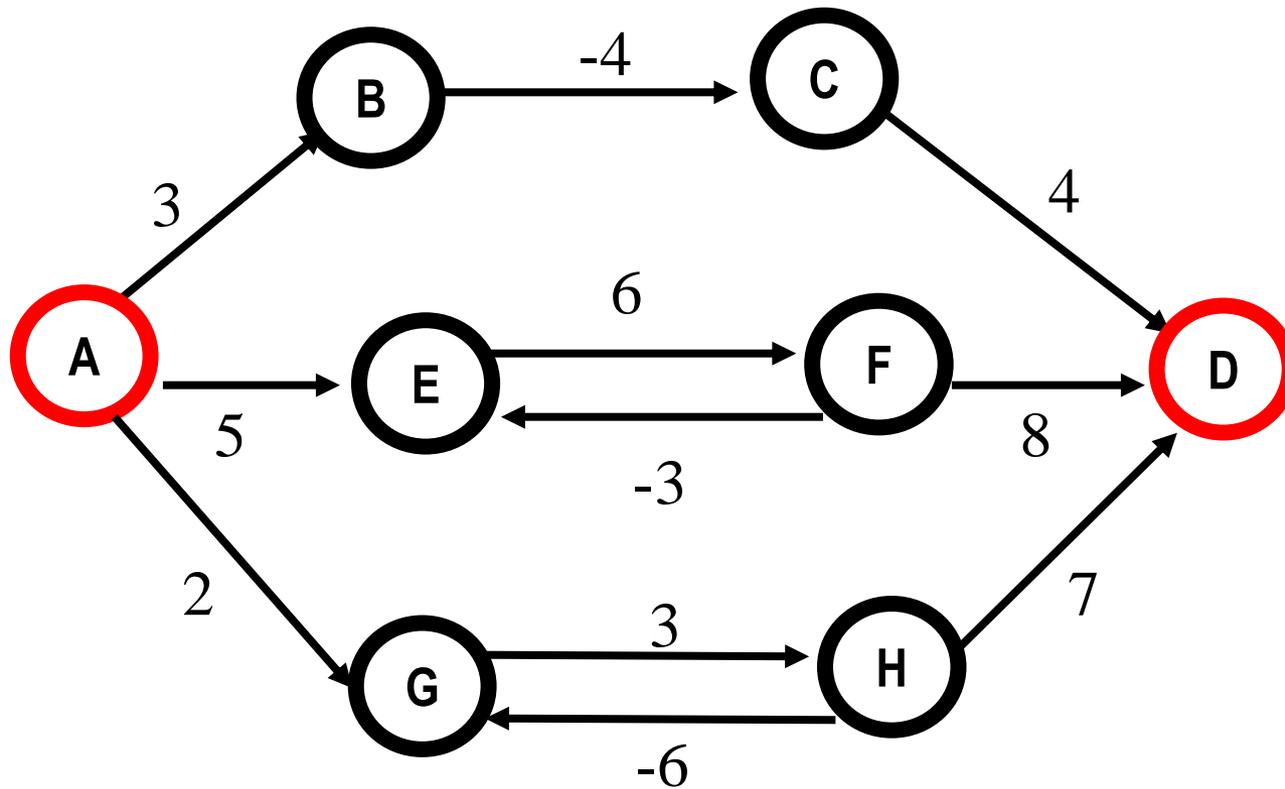
A shortest path from A to D ?

Example: Negative-Weight Edges + Nonnegative Cycles?



A shortest path from A to D ?

Example: Negative-Weight Edges + Negative Cycles?



A shortest path from A to D ?

The Shortest-Paths Problem with Negative Edges & Negative Cycles

- Is the shortest path problem defined for weighted directed graphs that contain negative edges but there is **no** negative cycle?
 - Yes
- Is the shortest path problem defined for weighted directed graphs that contain negative cost cycles?
 - No
- The shortest-path problem is greatly simplified when all edges have non-negative weights!

► QUIZ?

- Types of Shortest Path Problems?
 - SPSP
 - SSSP
 - SDSPP
 - APSP
- **TSP (Traveling Salesperson Problem)?**

Homework Assignment

▶ Homework Assignment?

- Compare **depth-first-search** with **breadth-first-search** of graphs.

the right majors, minors & concentrations
education?

for students' academic and career success
ing for many students - *Many students change their
uring college!*

e prediction of student success in MMC could
dual students
d their right MMC
hieve their academic goals

END

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Prof. Young Park

