

the right majors, minors & concentrations
education?
for students' academic and career success
ing for many students - *Many students change their
uring college!*

Branch-and-Bound (B&B)

e prediction of student success in MMC could
dual students
d their right MMC
hieve their academic goals

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177

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✓ Two Hard Optimization Problems So Far

The 0-1 Knapsack Problem

- **The 0-1 Knapsack problem** (optimization problem) solved so far.
 - **Via Greedy?**
 - No
 - **Via DP!**
 - $O(nW)$ or $O(2^n)$
 - **Any better?**
 - No one has ever found an algorithm whose worst-case time complexity is better than exponential!
 - Yet no one has proven that such an algorithm is not possible!

The Traveling Salesperson Problem

- **The Traveling Salesperson Problem** (optimization problem) solved so far.
 - **Via Greedy?**
 - No
 - **Via DP!**
 - $O(n^2 2^n)$
 - **Any better?**
 - No one has ever found an algorithm whose worst-case time complexity is better than exponential!
 - Yet no one has proven that such an algorithm is not possible!

Hard Optimization Problems

- Then, what?
- Still want to improve!
- How?
 - **Branch-and-Bound** (for *optimization problems*)

✓ Branch-and-Bound (B&B)

Branch-and-Bound (B&B)

- An approach/paradigm to designing algorithms!

Branch-and-Bound (B&B)

- The **branch-and-bound strategy** is very **similar to backtracking** in that a state space tree is used to solve a problem.
 - Does **not** limit us to **any particular way** of traversing the state space tree.
 - Is used only for **optimization problems**.
- Idea?

Bound

- A branch-and-bound algorithm computes a **number (Bound) at a node to determine whether the node is promising.**
 - The number is a **bound** on **the value of the solution that could be obtained by expanding beyond the node.**
 - If that bound is no better than the value of the best solution found so far, the node is **nonpromising.**
 - Otherwise it is promising.

Branch-and-Bound Search

- Does not limit us to any particular way of traversing the state space tree:
 - **Depth-first search**
 - Using **Stack**
 - **Breadth-first search**
 - Using **Queue**
 - **Best-first search**
 - Using **Priority Queue**

Best-First Search

- Besides using the bound to determine whether a node is promising,
- We can compare the bounds of promising nodes and visit the children of the one with the best bound.
- **Best-first search** with branch and bound pruning!
- Does not limit us to any particular way of traversing the state space tree!

Branch-and-Bound (B&B)

- The branch and bound algorithms **can be exponential or worse at the worst-case (like backtracking)**.
- The branch and bound algorithms can **often** arrive at an optimal solution **faster than backtracking's depth-first search**.
- The branch and bound algorithms are **efficient for many large instances**.

Branch-and-Bound (B&B)

- **An enhancement of Backtracking**
 - **Depth-first search**
 - Using **Stack**
 - **Breadth-first search**
 - Using **Queue**
 - **Best-first search**
 - Using **Priority Queue**
- **Applicable to optimization problems only**

▶ QUIZ?

- **Bound** is used for?
 - Ruling out certain nodes as “nonpromising” to prune the tree!
 - If a node’s bound is not better than the best solution seen so far, the node is non-promising.
 - Guiding the search through state-space!
 - Compare the bounds of promising nodes and visit the one with the best bound.

✓ Branch-and-Bound (B&B)-based Algorithms

Branch-and-Bound (B&B)-based Algorithms

1. The 0-1 Knapsack Problem via B&B
2. The Traveling Salesperson Problem via B&B

1. The 0-1 Knapsack problem via B&B

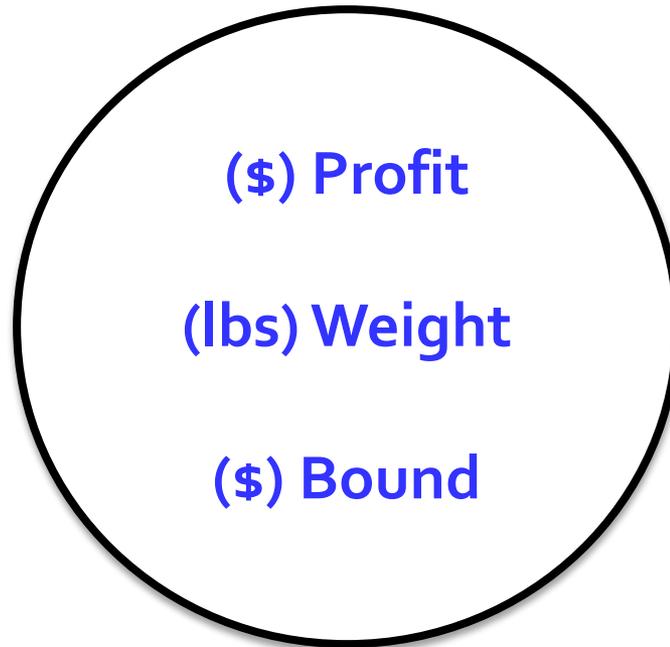
- The 0-1 Knapsack problem is an optimization problem.
- No polynomial time algorithm is known for the 0-1 Knapsack problem.
- Nobody has shown that a polynomial time algorithm is not possible for the 0-1 Knapsack problem.
- We apply B&B to the 0-1 Knapsack problem!

The 0-1 Knapsack Problem via B&B

- **Bound(v) (Potential profit upper bound)** is an **upper bound** on the **profit** we could achieve by expanding **beyond the node v!**
- **Pretending the Fractional Knapsack!**
- We will use **bound(v)** to determine whether **v** is non-promising!

The 0-1 Knapsack problem via B&B

- Each node consists of



The 0-1 Knapsack problem via B&B

- A node is **nonpromising** if
 1. The **total weight** thus far is greater than or equal to W .
 2. The **bound** (potential profit upper bound) is smaller than or equal to the value of **maxprofit**.

The 0-1 Knapsack problem via Breadth-First Search B&B

- **First**, a simple version called **breadth-first search with branch-and-bound pruning**.
 - Using a **queue**!

The 0-1 Knapsack problem via B&B (**Breadth-First Search** with Branch-and-Bound Pruning)

■ Example 6.1

- $n = 4$ & $W = 16$ lb
- Item₁ = \$40 & 2 lb
- Item₂ = \$30 & 5 lb
- Item₃ = \$50 & 10 lb
- Item₄ = \$10 & 5 lb

Sort by profit/weight

- **Item 1 & item 3 (Max profit= \$90)**

The 0-1 Knapsack Problem

State Space Tree

- $n = 4$ & $W = 16$ lb
 - Item1 = \$40 & 2 lb
 - Item2 = \$30 & 5 lb
 - Item3 = \$50 & 10 lb
 - Item4 = \$10 & 5 lb
- Sort by profit/weight
- 

$$\sum_{i=0, n} 2^i = 2^{(n+1)} - 1$$

Example A.3

Total 31 nodes

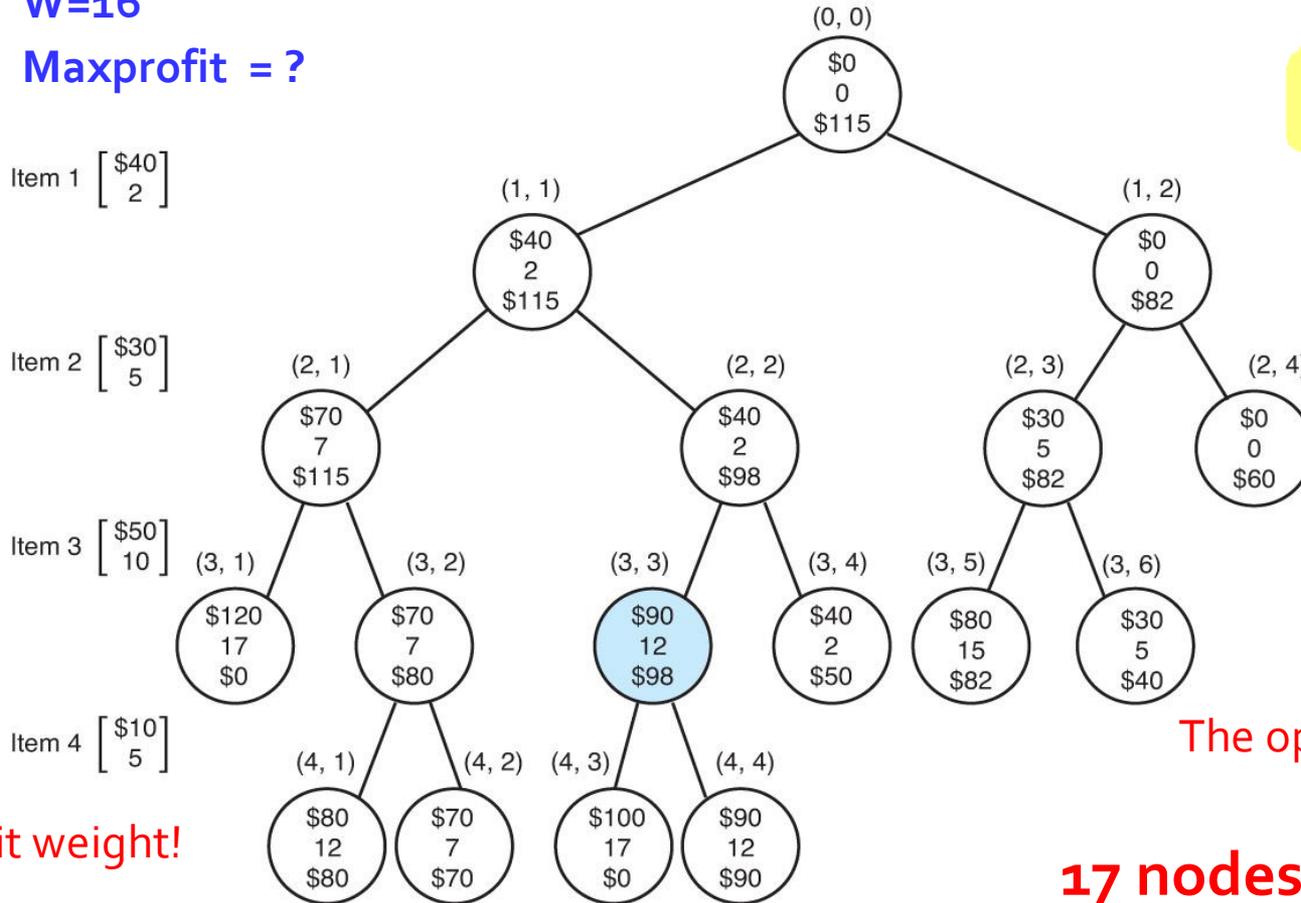
The 0-1 Knapsack problem via B&B (Breadth-First Search with Branch-and-Bound Pruning)

Figure 6.2



W=16
Maxprofit = ?

Total 31 nodes



Profit per unit weight!

17 nodes checked!

► QUIZ? The 0-1 Knapsack problem via B&B (Breadth-First Search)

- $W = 6$ lb
- Item₁ = \$10 & 1 lb
- Item₂ = \$18 & 2 lb
- Item₃ = \$32 & 4 lb
- Item₄ = \$14 & 2 lb

The 0-1 Knapsack problem via Best-First Search B&B

- **Second**, an improvement called **best-first search with branch-and-bound pruning**.
 - Compare the bounds of promising nodes and **visit the children of the one with the best (largest) bound**.
 - Using a **priority-queue!**
 - We **often arrive** at an optimal solution **more quickly** than if we simply proceeded blindly in a predetermined order.

The 0-1 Knapsack problem via B&B

(Best-First Search with Branch-and-Bound Pruning)

■ Example 6.2

- $n = 4$ & $W = 16$ lb
- Item₁ = \$40 & 2 lb
- Item₂ = \$30 & 5 lb
- Item₃ = \$50 & 10 lb
- Item₄ = \$10 & 5 lb

Sort by profit/weight

- **Item 1 & item 3 (Max profit = \$90)**

The 0-1 Knapsack Problem

State Space Tree

- $n = 4$ & $W = 16$ lb
 - Item1 = \$40 & 2 lb
 - Item2 = \$30 & 5 lb
 - Item3 = \$50 & 10 lb
 - Item4 = \$10 & 5 lb
- Sort by profit/weight
- 

$$\sum_{i=0, n} 2^i = 2^{(n+1)} - 1$$

Example A.3

Total 31 nodes

The 0-1 Knapsack problem via B&B

(Best-First Search with Branch-and-Bound Pruning)

Figure 6.3

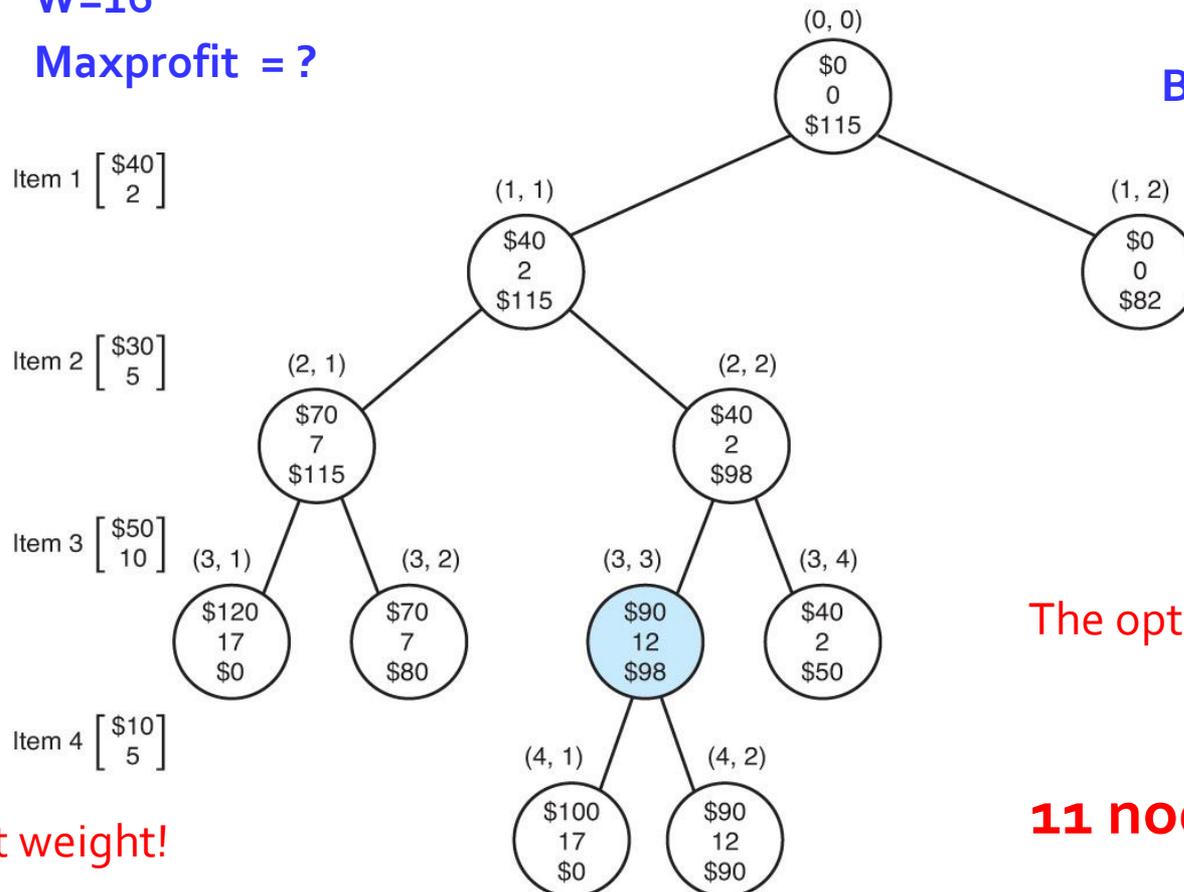


W=16

Maxprofit = ?

Total 31 nodes

Best-First Search



The optimal solution!

11 nodes checked!

Profit per unit weight!

► QUIZ? The 0-1 Knapsack problem via B&B (Best-First Search)

- $W = 6$ lb
- Item₁ = \$10 & 1 lb
- Item₂ = \$18 & 2 lb
- Item₃ = \$32 & 4 lb
- Item₄ = \$14 & 2 lb

The 0-1 Knapsack Problem

- **Recall:**
 - **Section 5.7**
 - **Using Backtracking for the 0-1 Knapsack Problem!**

The 0-1 Knapsack Problem via Backtracking

■ Example 5.6

- $n = 4$ & $W = 16$ lb
- Item₁ = \$40 & 2 lb
- Item₂ = \$30 & 5 lb
- Item₃ = \$50 & 10 lb
- Item₄ = \$10 & 5 lb

Sort by profit/weight

- **Item 1 & item 3 (Max profit = \$90)**

The 0-1 Knapsack Problem

State Space Tree

- $n = 4$ & $W = 16$ lb
 - Item1 = \$40 & 2 lb
 - Item2 = \$30 & 5 lb
 - Item3 = \$50 & 10 lb
 - Item4 = \$10 & 5 lb
- Sort by profit/weight
- 

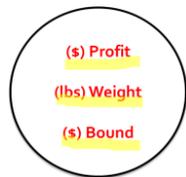
$$\sum_{i=0, n} 2^i = 2^{(n+1)} - 1$$

Example A.3

Total 31 nodes

The 0-1 Knapsack Problem – The State Space Tree Pruned via Backtracking

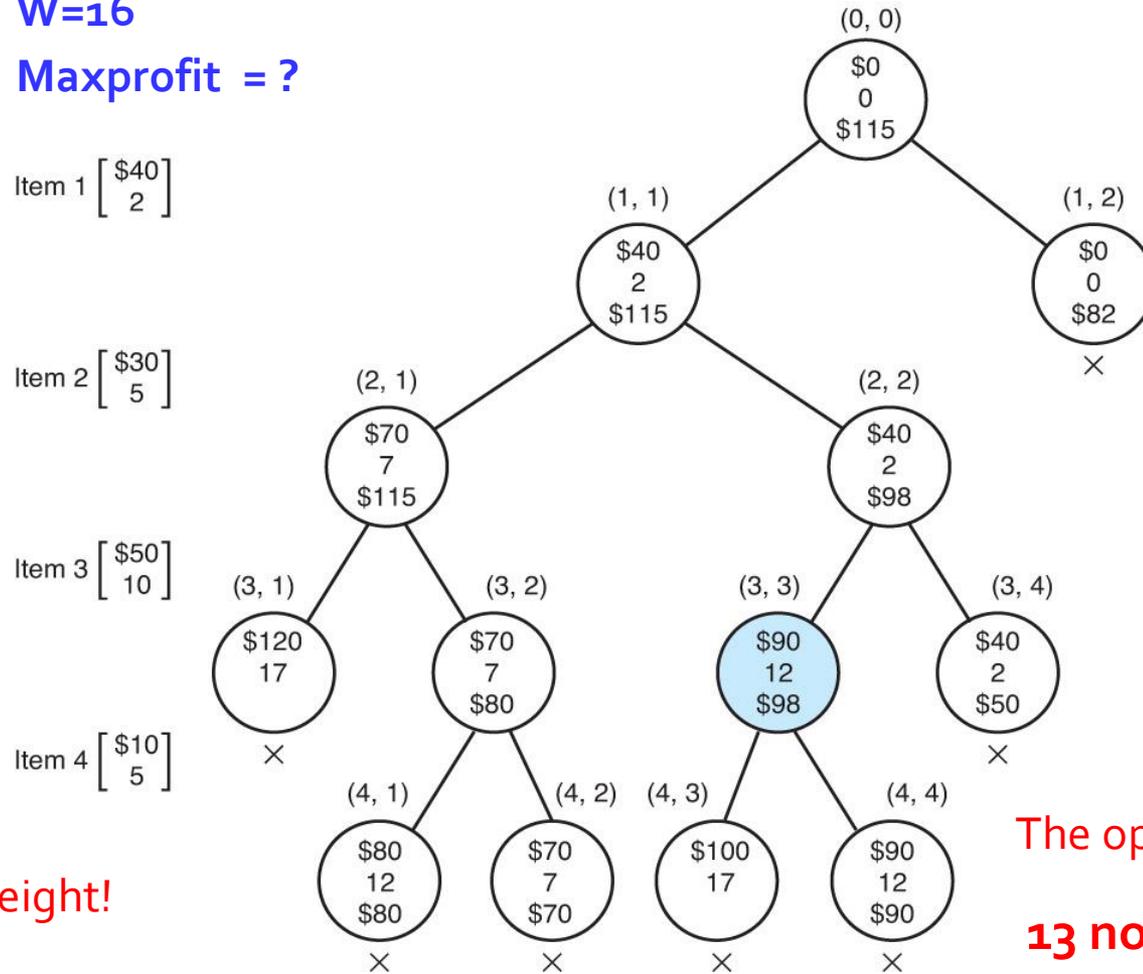
Figure 5.14



W=16
Maxprofit = ?

Total 31 nodes

DFS



Profit per unit weight!

The optimal solution!

13 nodes checked!

x= nonpromising node !

The 0-1 Knapsack Problem

- Actually,
 - Using **Backtracking** for the 0-1 Knapsack Problem!
 - This is a Depth-First Search with **branch-and-bound pruning**.
 - Using a **stack**!

2. The Traveling Salesperson Problem via B&B

- TSP is an optimization problem.
- No polynomial time algorithm is known for TSP.
- Nobody has shown that a polynomial time algorithm is not possible for TSP.
- **We can apply B&B to the TSP!**

▶ QUIZ?

- Compare **Brute-Force Approach** vs. **Backtracking** vs. **Branch-and-Bound**?

✓ Branch-and-Bound (B&B)-based Algorithms Summary

1. The 0-1 Knapsack Problem via B&B
2. The Traveling Salesperson Problem via B&B

Homework Assignment

► Homework Assignment?

- Draw the *pruned* state space tree via **Best-First Search B&B** and show the optimal solution of the **0-1 Knapsack Problem**: $W = 160$ lb & Item₁ = \$40 & 20 lb, Item₂ = \$30 & 50 lb, Item₃ = \$50 & 100 lb, Item₄ = \$10 & 50 lb.

✓ Textbook Readings

- Chapter 6:
 - 6.1

the right majors, minors & concentrations
education?

for students' academic and career success
ing for many students - *Many students change their
during college!*

the prediction of student success in MMC could
individual students
and their right MMC
achieve their academic goals

END

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1/7

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